

Chapter 2

Application Layer

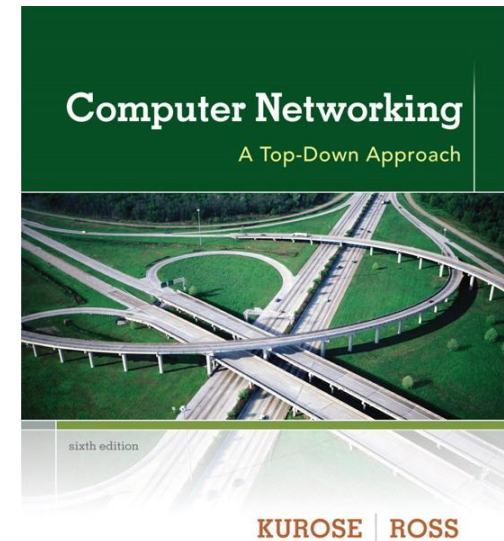
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*Computer
Networking: A
Top Down
Approach*
6th edition
Jim Kurose, Keith Ross
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Chapter 2: outline

2.1 principles of network applications

2.2 Web and HTTP

2.3 FTP

2.4 electronic mail

- SMTP, POP3, IMAP

2.5 DNS

2.6 P2P applications

2.7 socket programming with UDP and TCP

Some network apps

- ❖ e-mail
- ❖ web
- ❖ text messaging
- ❖ remote login
- ❖ P2P file sharing
- ❖ multi-user network games
- ❖ streaming stored video (YouTube, Hulu, Netflix)
- ❖ voice over IP (e.g., Skype)
- ❖ real-time video conferencing
- ❖ social networking
- ❖ search
- ❖ ...
- ❖ ...

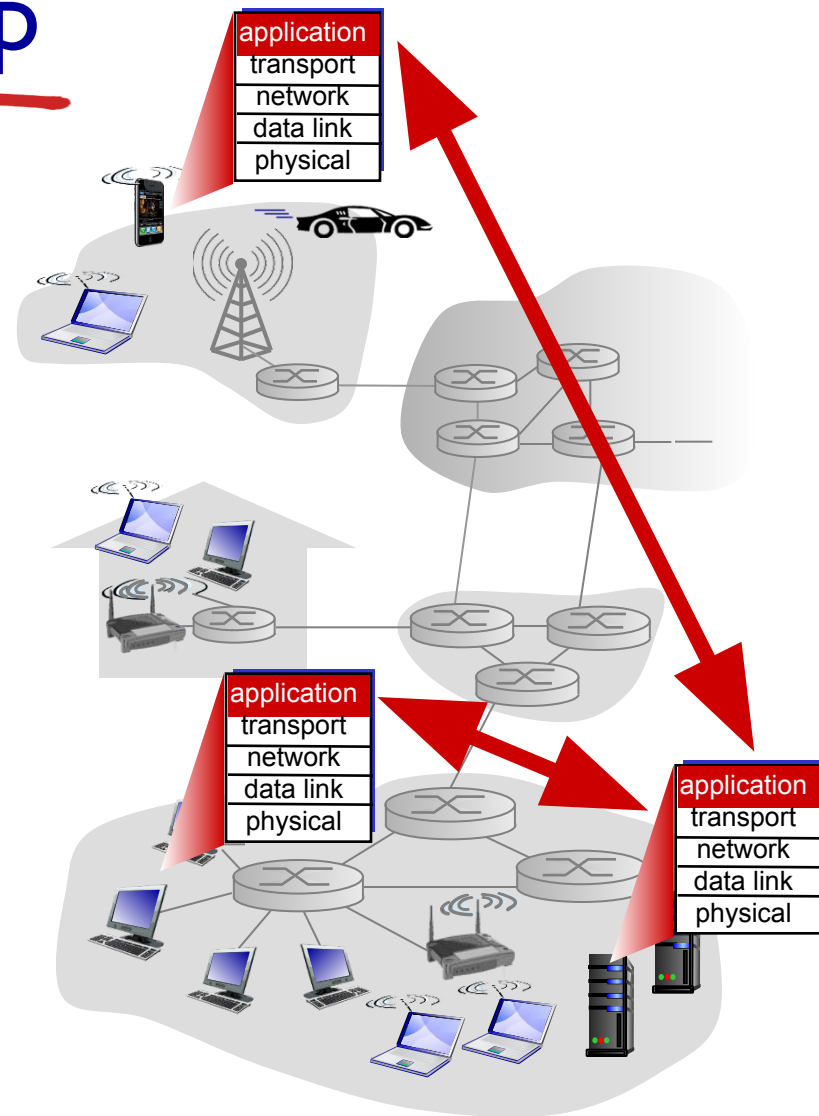
Creating a network app

write programs that:

- ❖ run on (different) *end systems*
- ❖ communicate over network
- ❖ e.g., web server software communicates with browser software

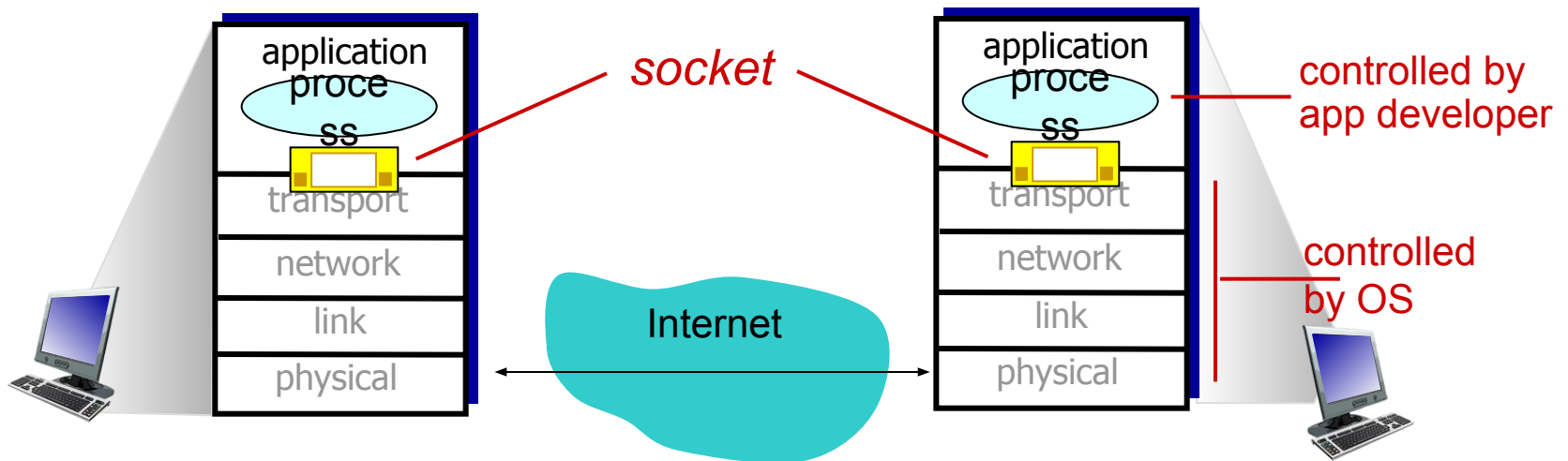
no need to write software for network-core devices

- ❖ network-core devices do not run user applications
- ❖ applications on end systems allows for rapid app development, propagation



Sockets

- ❖ process sends/receives messages to/from its **socket**
- ❖ socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



Addressing processes

- ❖ to receive messages, process must have *identifier*
- ❖ host device has unique 32-bit IP address
- ❖ Q: does IP address of host on which process runs suffice for identifying the process?
A: no, *many* processes can be running on same host
- ❖ *identifier* includes both **IP address** and **port numbers** associated with process on host.
- ❖ example port numbers:
 - HTTP server: 80
 - mail server: 25
- ❖ to send HTTP message to `gaia.cs.umass.edu` web server:
 - **IP address:** 128.119.245.12
 - **port number:** 80
- ❖ more shortly...

App-layer protocol defines

- ❖ **types of messages exchanged,**
 - e.g., request, response
- ❖ **message syntax:**
 - what fields in messages & how fields are delineated
- ❖ **message semantics**
 - meaning of information in fields
- ❖ **rules** for when and how processes send & respond to messages

open protocols:

- ❖ defined in RFCs
- ❖ allows for interoperability
- ❖ e.g., HTTP, SMTP

proprietary protocols:

- ❖ e.g., Skype

What transport service does an app need?

data integrity

- ❖ some apps (e.g., file transfer, web transactions) require 100% reliable data transfer
- ❖ other apps (e.g., audio) can tolerate some loss

timing

- ❖ some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

throughput

- ❖ some apps (e.g., multimedia) require minimum amount of throughput to be “effective”
- ❖ other apps (“elastic apps”) make use of whatever throughput they get

security

- ❖ encryption, data integrity, ...

Transport service requirements: common apps

application	data loss	throughput	time sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps video: 10kbps-5Mbps	yes, 100's msec
stored audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	few kbps up	yes, 100's msec
text messaging	no loss	elastic	yes and no

Internet transport protocols services

TCP service:

- ❖ *reliable transport* between sending and receiving process
- ❖ *flow control*: sender won't overwhelm receiver
- ❖ *congestion control*: throttle sender when network overloaded
- ❖ *does not provide*: timing, minimum throughput guarantee, security
- ❖ *connection-oriented*: setup required between client and server processes

UDP service:

- ❖ *unreliable data transfer* between sending and receiving process
 - ❖ *does not provide*: reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup,
- Q: why bother? Why is there a UDP?

Internet apps: application, transport protocols

	application	application layer protocol	underlying transport protocol
	e-mail	SMTP [RFC 2821]	TCP
remote terminal access		Telnet [RFC 854]	TCP
	Web	HTTP [RFC 2616]	TCP
	file transfer	FTP [RFC 959]	TCP
streaming multimedia		HTTP (e.g., YouTube), RTP [RFC 1889]	TCP or UDP
Internet telephony		SIP, RTP, proprietary (e.g., Skype)	TCP or UDP

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2.1 principles of network applications

- app architectures
- app requirements

2.2 Web and HTTP

2.3 FTP

2.4 electronic mail

- SMTP, POP3, IMAP

2.5 DNS

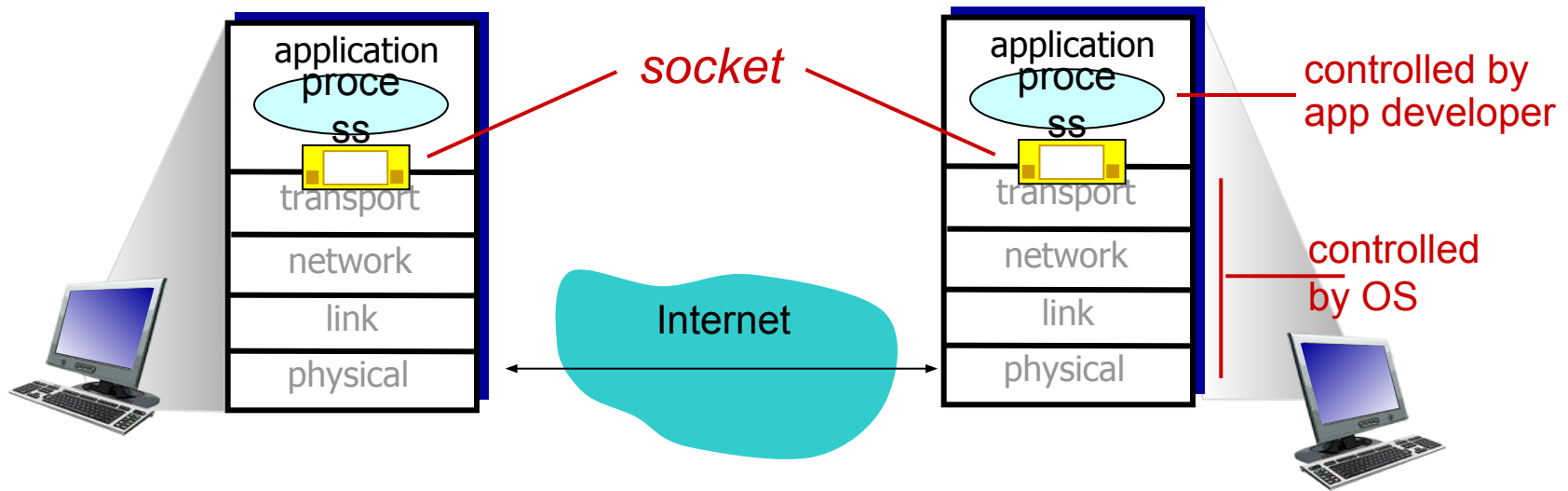
2.6 P2P applications

2.7 socket programming with UDP and TCP

Socket programming

goal: learn how to build client/server applications that communicate using sockets

socket: door between application process and end-end-transport protocol



Socket programming

Two socket types for two transport services:

- **UDP:** unreliable datagram
- **TCP:** reliable, byte stream-oriented

Application Example:

1. Client reads a line of characters (data) from its keyboard and sends the data to the server.
2. The server receives the data and converts characters to uppercase.
3. The server sends the modified data to the client.
4. The client receives the modified data and displays the line on its screen.

Socket programming *with UDP*

UDP: no “connection” between client & server

- ❖ no handshaking before sending data
- ❖ sender explicitly attaches IP destination address and port # to each packet
- ❖ rcvr extracts sender IP address and port# from received packet

UDP: transmitted data may be lost or received out-of-order

Application viewpoint:

- ❖ UDP provides *unreliable* transfer of groups of bytes (“datagrams”) between client and server

Client/server socket interaction: UDP

server (running on *serverIP*)

create socket, port= x:
`serverSocket =
socket(AF_INET,SOCK_DGRAM)`

read datagram from
`serverSocket`

write reply to
`serverSocket`
specifying
client address,
port number

client

create socket:
`clientSocket =
socket(AF_INET,SOCK_DGRAM)`

Create datagram with server IP and
port=x; send datagram via
`clientSocket`

read datagram from
`clientSocket`

close
`clientSocket`

Example app: UDP client

Python UDPClient

include Python's socket
library

→ from socket import *

serverName = 'hostname'

serverPort = 12000

create UDP socket for
server

→ clientSocket = socket(socket.AF_INET,

socket.SOCK_DGRAM)

get user keyboard
input

→ message = raw_input('Input lowercase sentence:')

Attach server name, port to
message; send into socket

→ clientSocket.sendto(message,(serverName, serverPort))

read reply characters from
socket into string

modifiedMessage, serverAddress =

print out received string
and close socket

→ clientSocket.recvfrom(2048)

print modifiedMessage

clientSocket.close()

Example app: UDP server

Python UDPServer

```
from socket import *  
  
serverPort = 12000  
  
create UDP socket → serverSocket = socket(AF_INET, SOCK_DGRAM)  
bind socket to local port  
number 12000 → serverSocket.bind(("", serverPort))  
  
loop forever → print "The server is ready to receive"  
  
Read from UDP socket into  
message, getting client's  
address (client IP and port) → while 1:  
    message, clientAddress = serverSocket.recvfrom(2048)  
    send upper case string  
    back to this client → modifiedMessage = message.upper()  
    serverSocket.sendto(modifiedMessage, clientAddress)
```

Socket programming *with TCP*

client must contact server

- ❖ server process must first be running
- ❖ server must have created socket (door) that welcomes client's contact

client contacts server by:

- ❖ Creating TCP socket, specifying IP address, port number of server process
- ❖ *when client creates socket:* client TCP establishes connection to server TCP

- ❖ when contacted by client, *server TCP creates new socket* for server process to communicate with that particular client
 - allows server to talk with multiple clients
 - source port numbers used to distinguish clients (more in Chap 3)

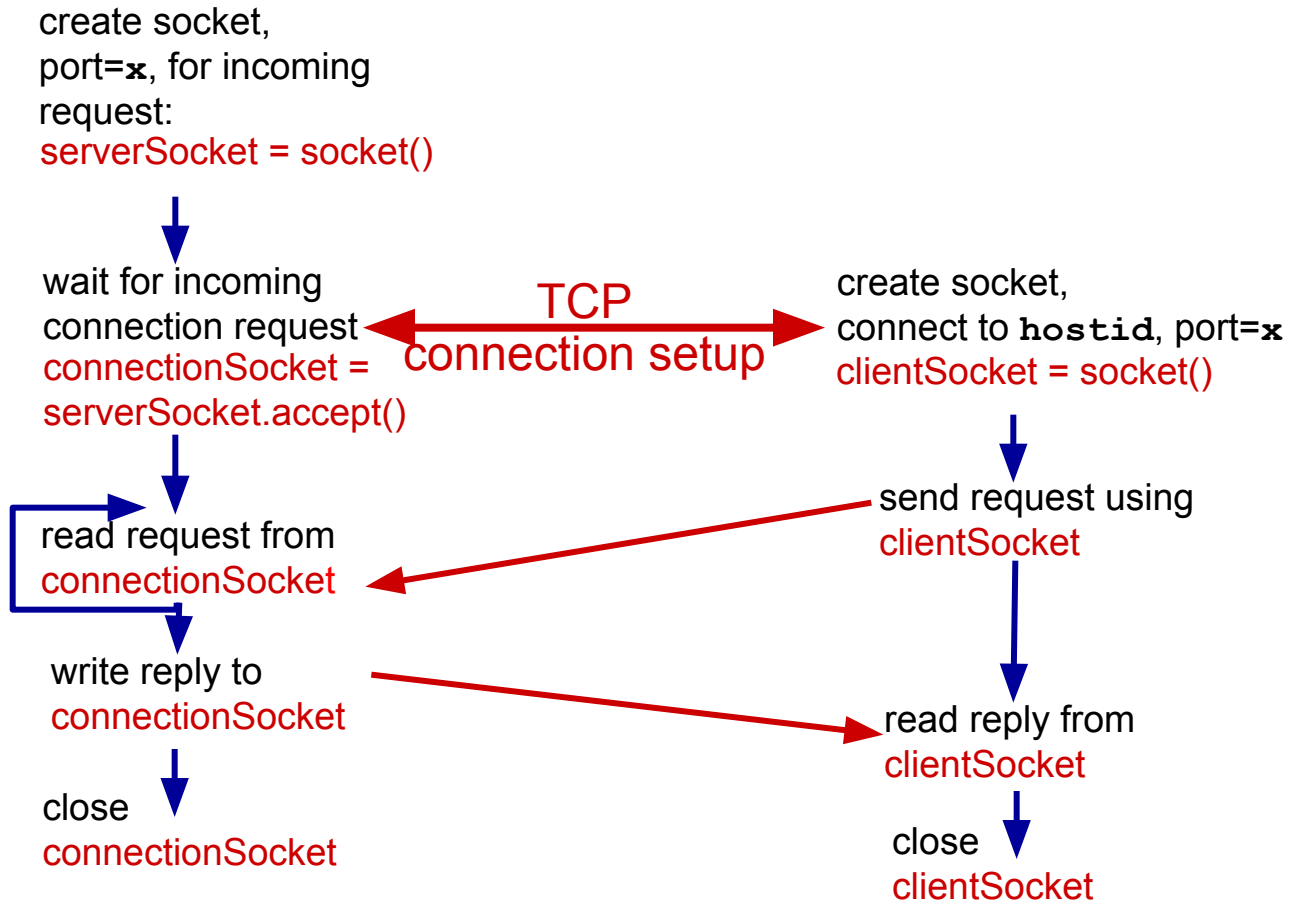
application viewpoint:

TCP provides reliable, in-order byte-stream transfer (“pipe”) between client and server

Client/server socket interaction: TCP

server (running on `hostid`)

client



Example app:TCP client

Python TCPClient

```
from socket import *  
  
serverName = 'servername'  
  
serverPort = 12000  
  
clientSocket = socket(AF_INET, SOCK_STREAM)  
clientSocket.connect((serverName,serverPort))  
  
sentence = raw_input('Input lowercase sentence:')  
clientSocket.send(sentence)  
  
modifiedSentence = clientSocket.recv(1024)  
  
print 'From Server:', modifiedSentence  
  
clientSocket.close()
```

create TCP socket for
server, remote port 12000



No need to attach server
name, port



Example app: TCP server

Python TCPServer

create TCP welcoming
socket



server begins listening for
incoming TCP requests



loop forever



server waits on accept()
for incoming requests, new
socket created on return



read bytes from socket (but
not address as in UDP)



close connection to this
client (but *not* welcoming
socket)



```
from socket import *
serverPort = 12000
serverSocket = socket(AF_INET,SOCK_STREAM)
serverSocket.bind(('',serverPort))
serverSocket.listen(1)
print 'The server is ready to receive'
while 1:
    connectionSocket, addr = serverSocket.accept()
    sentence = connectionSocket.recv(1024)
    capitalizedSentence = sentence.upper()
    connectionSocket.send(capitalizedSentence)
    connectionSocket.close()
```