



# Decompiler internals: microcode

Hex-Rays  
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# Presentation Outline

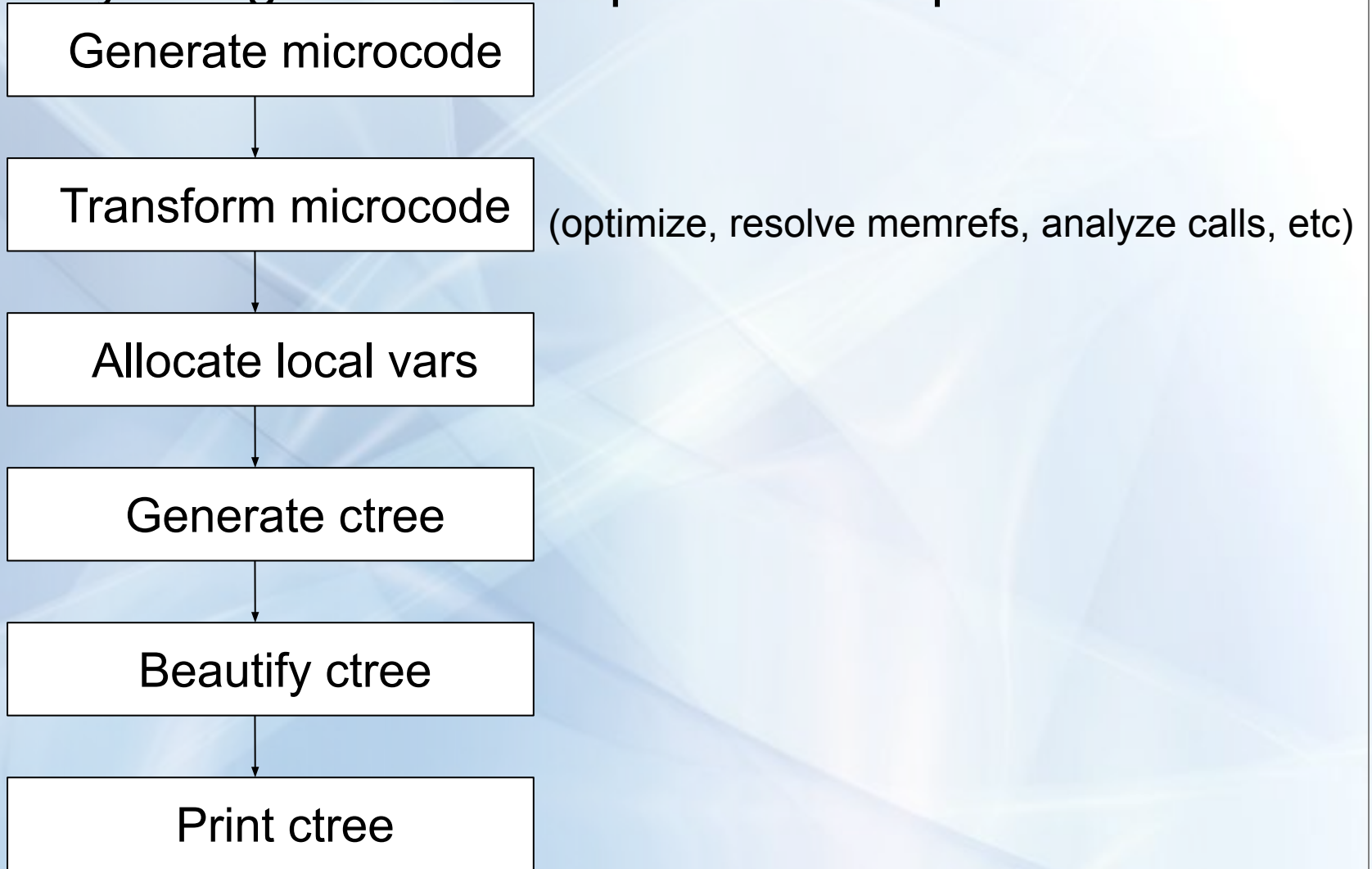
- Decompiler architecture
  - Overview of the microcode
  - Opcodes and operands
  - Stack and registers
  - Data flow analysis, aliasability
  - Microcode availability
  - Your feedback
- 
- Online copy of this presentation is available at  
<http://www.hex-rays.com/products/ida/support/ppt/recon2018.ppt>

# Hex-Rays Decompiler

- Interactive, fast, robust, and programmable decompiler
- Can handle x86, x64, ARM, ARM64, PowerPC
- Runs on top of IDA Pro
- Has been evolving for more than 10 years
- Internals were not really published
- Namely, the intermediate language

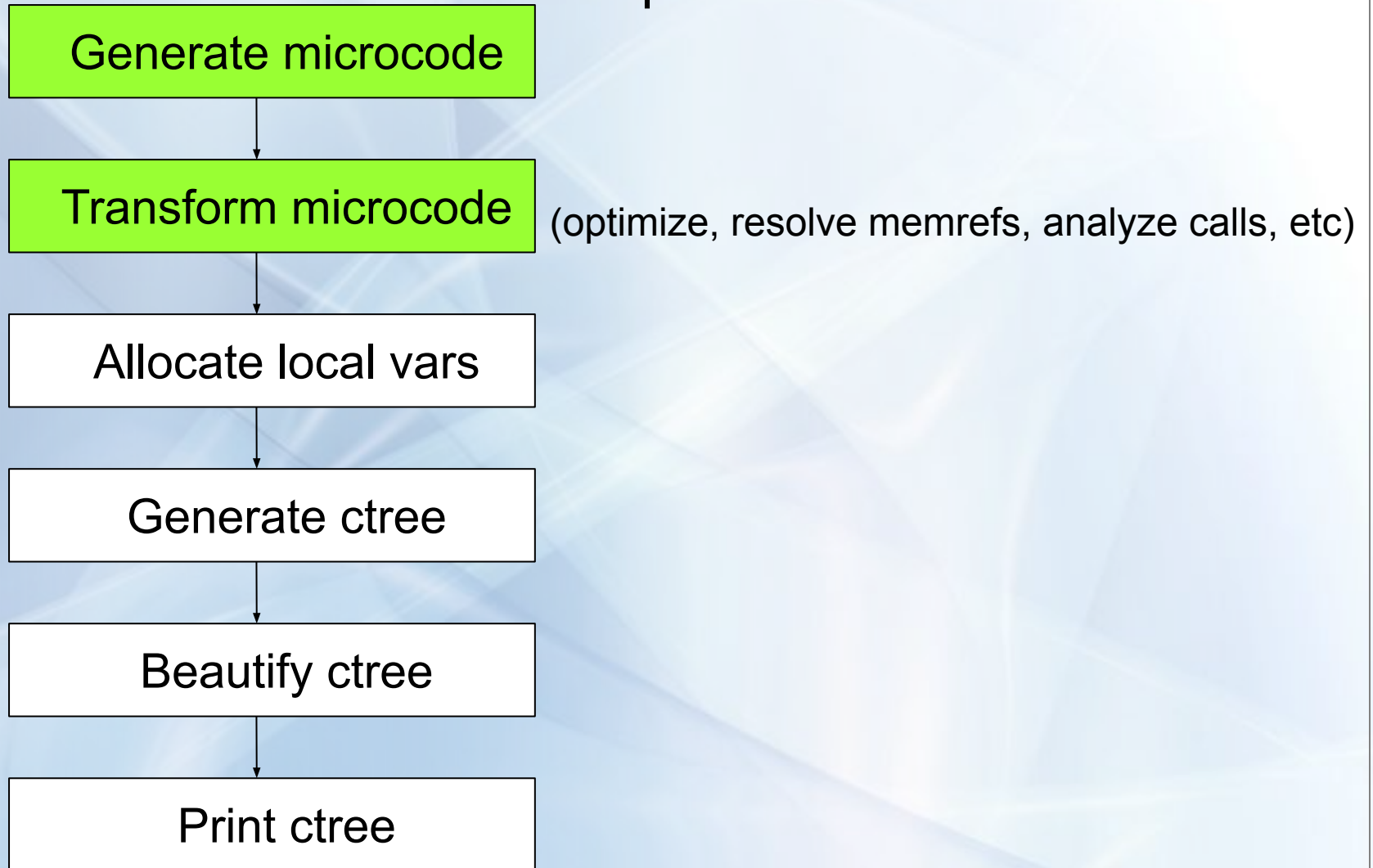
# Decompiler architecture

- It uses very straightforward sequence of steps:



# Decompiler architecture

- We will focus on the first two steps:



## Why microcode?

- It helps to get rid of the complexity of processor instructions
- Also we get rid of processor idiosyncrasies. Examples:
  - x86: segment registers, fpu stack
  - ARM: thumb mode addresses
  - PowerPC: multiple copies of CF register (and other condition registers)
  - MIPS: delay slots
  - Sparc: stack windows
- It makes the decompiler portable. We “just” need to replace the microcode generator
- Writing a decompiler without an intermediate language looks like waste of time

## Is implementing an IR difficult?

- Your call :)
- How many IR languages to you know?

## Why not use an existing IR?

- There are tons of other intermediate languages: LLVM, REIL, Binary Ninja's ILs, RetDec's IL, etc.
- Yes, we could use something
- But I started to work on the microcode when none of the above languages existed
- This is the main reason why we use our own IR

```
mov.d  EAX,, T0
ldc.d  #5,, T1
mkcadd.d T0, T1, CF
mkoadd.d T0, T1, CF
add.d  T0, T1, TT
setz.d  TT,, ZF
sets.d  TT,, ZF
mov.d  TT,, EAX
```

(this is how it looked like in 1999)



## A long evolution

- I started to work on the microcode in 1998 or earlier
- The name is nothing fancy but reflects the nature of it
- Some design decisions turned out to be bad (and some of them are already very difficult to fix)
- For example, the notion of virtual stack registers
- We will fix it, though. Just takes time
- Even today we modify our microcode when necessary
- For example, I reshuffled the instruction opcodes for this talk...

# Design highlights

- Simplicity:
  - No processor specific stuff
  - One microinstruction does one thing
  - Small number of instructions (only 45 in 1999, now 72)
  - Simple instruction operands (register, number, memory)
  - Consider only compiler generated code
- Discard things we do not care about:
  - Instruction timing (anyway it is a lost battle)
  - Instruction order (exceptions are a problem!)
  - Order of memory accesses (later we added logic to preserve indirect memory accesses)
  - Handcrafted code

## Generated microcode

- Initially the microcode looks like RISC code:
  - Memory loads and stores are done using dedicated microinstructions
  - The desired operation is performed on registers
  - Microinstructions have no side effects
  - Each output register is initialized by a separate microinstruction
- It is very verbose. Example:

```
004014FB    mov    eax, [ebx+4]
004014FE    mov    dl, [eax+1]
00401501    sub    dl, 61h ; 'a'
00401504    jz     short loc_401517
```

# Initial microcode: very verbose

```
2. 0 mov    ebx.4, eoff.4          ; 4014FB u=ebx.4      d=eoff.4
2. 1 mov    ds.2, seg.2            ; 4014FB u=ds.2      d=seg.2
2. 2 add    eoff.4, #4.4, eoff.4   ; 4014FB u=eoff.4    d=eoff.4
2. 3 ldx    seg.2, eoff.4, et1.4   ; 4014FB u=eoff.4,seg.2,
                                   ; (STACK,GLBMEM) d=et1.4
2. 4 mov    et1.4, eax.4           ; 4014FB u=et1.4     d=eax.4
2. 5 mov    eax.4, eoff.4          ; 4014FE u=eax.4     d=eoff.4
2. 6 mov    ds.2, seg.2            ; 4014FE u=ds.2     d=seg.2
2. 7 add    eoff.4, #1.4, eoff.4   ; 4014FE u=eoff.4    d=eoff.4
2. 8 ldx    seg.2, eoff.4, t1.1    ; 4014FE u=eoff.4,seg.2,
                                   ; (STACK,GLBMEM) d=t1.1
2. 9 mov    t1.1, dl.1             ; 4014FE u=t1.1     d=dl.1
2.10 mov    #0x61.1, t1.1         ; 401501 u=          d=t1.1
2.11 setb   dl.1, t1.1, cf.1       ; 401501 u=dl.1,t1.1 d=cf.1
2.12 seto   dl.1, t1.1, of.1       ; 401501 u=dl.1,t1.1 d=of.1
2.13 sub    dl.1, t1.1, dl.1       ; 401501 u=dl.1,t1.1 d=dl.1
2.14 setz   dl.1, #0.1, zf.1       ; 401501 u=dl.1     d=zf.1
2.15 setp   dl.1, #0.1, pf.1       ; 401501 u=dl.1     d=pf.1
2.16 sets   dl.1, sf.1             ; 401501 u=dl.1     d=sf.1
2.17 mov    cs.2, seg.2            ; 401504 u=cs.2     d=seg.2
2.18 mov    #0x401517.4, eoff.4    ; 401504 u=          d=eoff.4
2.19 jcnd   zf.1, $loc_401517     ; 401504 u=zf.1
```

## The first optimization pass

```
2. 0  ldx    ds.2, (ebx.4+#4.4), eax.4 ; 4014FB u=ebx.4,ds.2,
                                           ;(STACK,GLBMEM) d=eax.4
2. 1  ldx    ds.2, (eax.4+#1.4), dl.1 ; 4014FE u=eax.4,ds.2,
                                           ;(STACK,GLBMEM) d=dl.1
2. 2  setb   dl.1, #0x61.1, cf.1      ; 401501 u=dl.1      d=cf.1
2. 3  seto   dl.1, #0x61.1, of.1      ; 401501 u=dl.1      d=of.1
2. 4  sub    dl.1, #0x61.1, dl.1      ; 401501 u=dl.1      d=dl.1
2. 5  setz   dl.1, #0.1, zf.1         ; 401501 u=dl.1      d=zf.1
2. 6  setp   dl.1, #0.1, pf.1         ; 401501 u=dl.1      d=pf.1
2. 7  sets   dl.1, sf.1               ; 401501 u=dl.1      d=sf.1
2. 8  jcnd   zf.1, $loc_401517        ; 401504 u=zf.1
```

- Only 8 microinstructions
- Some intermediate registers disappeared
- Sub-instructions appeared
- Still too noisy and verbose

## Further microcode transformations

```
2. 1 ldx    ds.2{3}, ([ds.2{3}:(ebx.4+#4.4)].4+#1.4), dl.1{5} ; 4014FE
           ; u=ebx.4,ds.2,(GLBLow,sp+20..,GLBHIGH) d=dl.1
2. 2 sub    dl.1{5}, #0x61.1, dl.1{6} ; 401501 u=dl.1      d=dl.1
2. 3 jz     dl.1{6}, #0.1, @7         ; 401504 u=dl.1
```

And the final code is:

```
2. 0 jz     [ds.2{4}:([ds.2{4}:(ebx.4{8}+#4.4){7}].4{6}+#1.4){5}].1{3},
           #0x61.1,
           @7
           ; 401504 u=ebx.4,ds.2,(GLBLow,GLBHIGH)
```

This code is ready to be translated to ctree.  
(numbers in curly braces are value numbers)

The output will look like this:

```
if ( argv[1][1] == 'a' )
    ...
```

## Minor details

- Reading microcode is not easy (but hey, it was not designed for that! :)
- All operand sizes are spelled out explicitly
- The initial microcode is very simple (RISC like)
- As we transform microcode, nested subinstructions may appear
- We implemented the translation from processor instructions to microinstructions in plain C++
- We do not use automatic code generators or machine descriptions to generate them. Anyway there are too many processor specific details to make them feasible

## Opcodes: constants and move

- Copy from (l) to (d)estination
- Operand sizes must match

```
ldc  l,    d    // load constant
mov  l,    d    // move
```



## Opcodes: changing operand size

- Copy from (l) to (d)estination
- Operand sizes must differ
- Since real world programs work with partial registers (like al, ah), we absolutely need low/high

```
xds l,    d    // extend (signed)
xdu l,    d    // extend (unsigned)
low l,    d    // take low part
high l,   d    // take high part
```

## Opcodes: load and store

- {sel, off} is a segment:offset pair
- Usually seg is ds or cs; for processors with flat memory it is ignored
- 'off' is the most interesting part, it is a memory address

```
stx  l, sel, off // store value to memory
ldx  sel, off, d // load value from memory
```

Example:

```
ldx  ds.2, (ebx.4+#4.4), eax.4
stx  #0x2E.1, ds.2, eax.4
```

## Opcodes: comparisons

- Compare (l)left against (r)right
- The result is stored into (d)estination, a bit register like CF,ZF,SF,...

```
sets  l,    d // sign
setp  l, r, d // unordered/parity
setnz l, r, d // not equal
setz  l, r, d // equal
setae l, r, d // above or equal
setb  l, r, d // below
seta  l, r, d // above
setbe l, r, d // below or equal
setg  l, r, d // greater
setge l, r, d // greater or equal
setl  l, r, d // less
setle l, r, d // less or equal
seto  l, r, d // overflow of (l-r)
```

# Opcodes: arithmetic and bitwise operations

- Operand sizes must be the same
- The result is stored into (d)estination

```
neg  l,    d    // -l    -> d
lnot l,    d    // !l    -> d
bnot l,    d    // ~l    -> d
add  l, r,  d    // l + r -> d
sub  l, r,  d    // l - r -> d
mul  l, r,  d    // l * r -> d
udiv l, r,  d    // l / r -> d
sdiv l, r,  d    // l / r -> d
umod l, r,  d    // l % r -> d
smod l, r,  d    // l % r -> d
or   l, r,  d    // bitwise or
and  l, r,  d    // bitwise and
xor  l, r,  d    // bitwise xor
```

## Opcodes: shifts (and rotations?)

- Shift (l)eft by the amount specified in (r)ight
- The result is stored into (d)estination
- Initially our microcode had rotation operations but they turned out to be useless because they can not be nicely represented in C

```
shl  l, r, d    // shift logical left
shr  l, r, d    // shift logical right
sar  l, r, d    // shift arithmetic right
```

## Opcodes: condition codes

- Perform the operation on (l)left and (r)ight
- Generate carry or overflow bits
- Store CF or OF into (d)estination
- We need these instructions to precisely track carry and overflow bits
- Normally these instructions get eliminated during microcode transformations

```
cfadd l, r, d    // carry    of (l+r)
ofadd l, r, d    // overflow of (l+r)
cfshl l, r, d    // carry    of (l<<r)
cfshr l, r, d    // carry    of (l>>r)
```

## Opcodes: unconditional flow control

- Initially calls have only the callee address
- The decompiler retrieves the callee prototype from the database or tries to guess it
- After that the 'd' operand contains all information about the call, including the function prototype and actual arguments

```
ijmp {sel, off}    // indirect jmp
goto l             // unconditional jmp
call l            d // direct call
icall {sel, off} d // indirect call
ret               // return
```

```
call  $__org_fprintf <....:
      "FILE *" &($stdout).4,
      "const char *" &($aArIllegalSwitc).4,
      _DWORD xds.4([ds.2:([ds.2:(ebx.4+#4.4)].4+#1.4)].1)>.0
```

## Opcodes: conditional jumps

- Compare (l)eft against (r)ight and jump to (d)estination if the condition holds
- Jtbl is used to represent 'switch' idioms

```
jcnd  l,    d    //  
jnz   l, r, d    // ZF=0      Not Equal  
jz    l, r, d    // ZF=1      Equal  
jae   l, r, d    // CF=0      Above or Equal  
jb    l, r, d    // CF=1      Below  
ja    l, r, d    // CF=0 & ZF=0 Above  
jbe   l, r, d    // CF=1 | ZF=1 Below or Equal  
jg    l, r, d    // SF=OF & ZF=0 Greater  
jge   l, r, d    // SF=OF      Greater or Equal  
jl    l, r, d    // SF!=OF     Less  
jle   l, r, d    // SF!=OF | ZF=1 Less or Equal  
jtbl  l, cases  // Table jump
```



## Opcodes: floating point operations

- Basically we have conversions and a few arithmetic operations
- There is little we can do with these operations, they are not really optimizable
- Other fp operations use helper functions (e.g. sqrt)

```
f2i   l,    d    // int(l) => d; convert fp -> int, any size
f2u   l,    d    // uint(l)=> d; convert fp -> uint, any size
i2f   l,    d    // fp(l)  => d; convert int -> fp, any size
i2f   l,    d    // fp(l)  => d; convert uint-> fp, any size
f2f   l,    d    // l      => d; change fp precision
fneg  l,    d    // -l     => d; change sign
fadd  l, r,  d    // l + r  => d; add
fsub  l, r,  d    // l - r  => d; subtract
fmul  l, r,  d    // l * r  => d; multiply
fdiv  l, r,  d    // l / r  => d; divide
```

## Opcodes: miscellaneous

- Some operations can not be expressed in microcode
- If possible, we use intrinsic calls for them (e.g. sqrtpd)
- If no intrinsic call exists, we use “ext” for them and only try to keep track of data dependencies (e.g. “aam”)
- “und” is used when a register is spoiled in a way that we can not predict or describe (e.g. ZF after mul)

```
nop                // no operation
und                d    // undefine
ext  l, r, d       // external insn
push l
pop               d
```

## More opcodes?

- We quickly reviewed all 72 instructions
- Probably we should extend microcode
- Ternary operator?
- Post-increment and post-decrement?
- All this requires more research

# Operands!

- As everyone else, initially we had only:
  - constant integer numbers
  - registers
- Life was simple and easy in the good old days!
- Alas, the reality is more diverse. We quickly added:
  - stack variables
  - global variables
  - address of an operand
  - list of cases (for switches)
  - result of another instruction
  - helper functions
  - call arguments
  - string and floating point constants

## Register operands

- The microcode engine provides unlimited (in theory) number of microregisters
- Processor registers are mapped to microregisters:
  - `eax` => microregisters (mreg) 8, 9, 10, 11
  - `al` => mreg 8
  - `ah` => mreg 9
- Usually there are more microregisters than the processor registers. We allocate them as needed when generating microcode
- Examples:

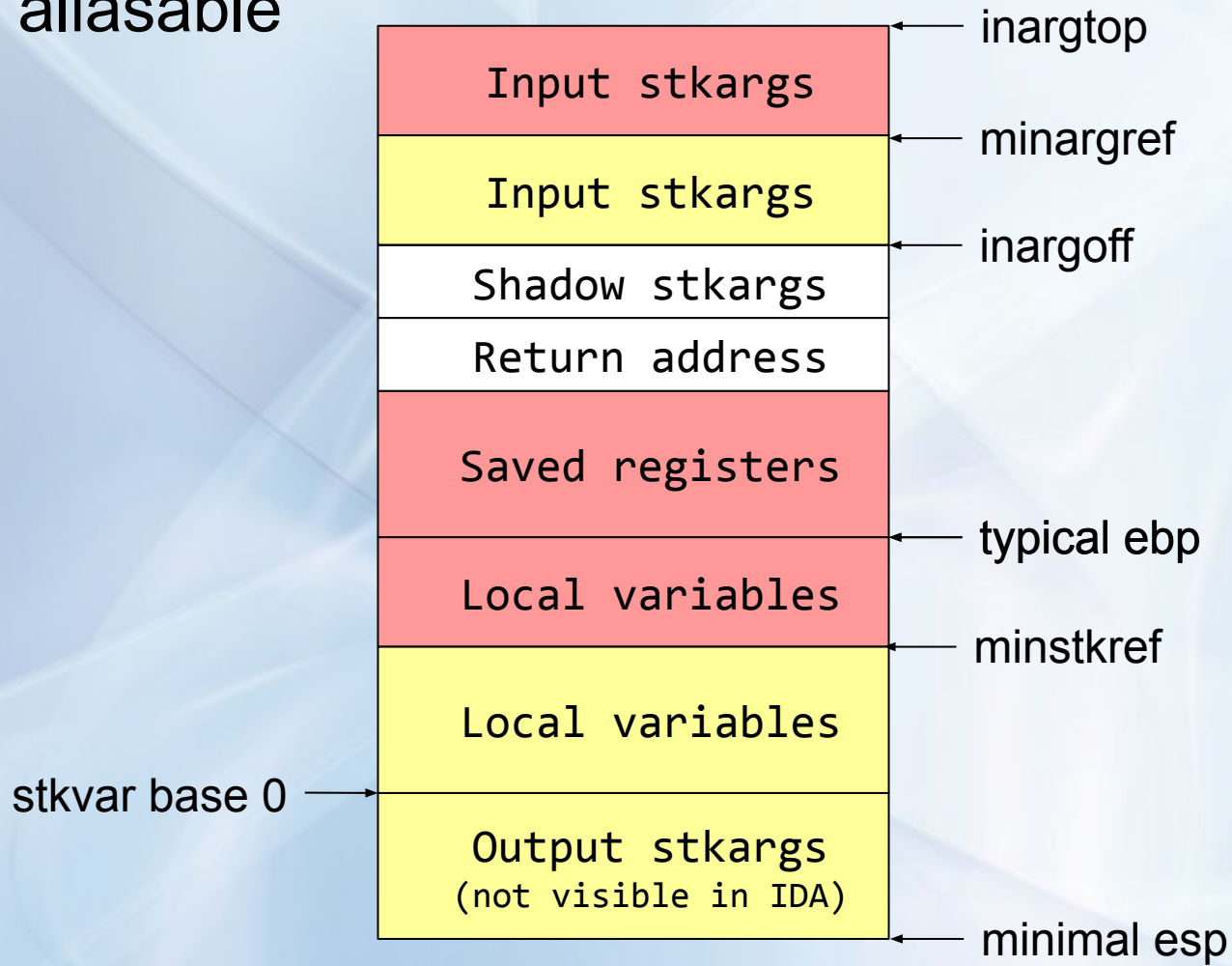
```
eax.4  
rsi.8  
ST00_04.4
```

## Stack as microregisters

- I was reluctant to introduce a new operand type for stack variables and decided to map the stack frame to microregisters
- Like, the stack frame is mapped to the microregister #100 and higher
- A bright idea? Nope!
- Very soon I realized that we have to handle indirect references to the stack frame
- Not really possible with microregisters
- But there was so much code relying on this concept that we still have it
- Laziness pays off now and in the future (negatively)

# Stack as viewed by the decompiler

- Yellow part is mapped to microregisters
- Red is aliasable



## More operand types!

- 64-bit values are represented as pairs of registers
- Usually it is a standard pair like edx:eax
- Compilers get better and nowadays use any registers as a pair; or even pair a stack location with a register: sp+4:esi
- We ended up with a new operand type:
  - operand pair
- It consists of low and high halves
- They can be located anywhere (stack, registers, glbmem)



## Scattered operands

- The nightmare has just begun, in fact
- Modern compilers use very intricate rules to pass structs and unions by value to and from the called functions
- A register like RDI may contain multiple structure fields
- Some structure fields may be passed on the stack
- Some in the floating registers
- Some in general registers (unaligned wrt register start)
- We had no other choice but to add
  - scattered operands

that can represent all the above

# A simple scattered return value

- A function that returns a struct in rax:

```
struct div_t { int quot; int rem; };  
div_t div(int numer, int denom);
```

- Assembler code:

```
mov     edi, esi  
mov     esi, 1000  
call   _div  
movsxd rdx, eax  
sar    rax, 20h  
add    [rbx], rdx  
imul   eax, 1000  
cdqe  
add    rax, [rbx+8]
```

## A simple scattered return value

- ...and the output is:

```
v2 = div(a2, 1000);  
*a1 += v2.quot;  
result = a1[1] + 1000 * v2.rem;
```

- Our decompiler managed to represent things nicely!
- Similar or more complex situations exist for all 64-bit processors
- Support for scattered operands is not complete yet but we constantly improve it

## More detailed look at microcode transformations

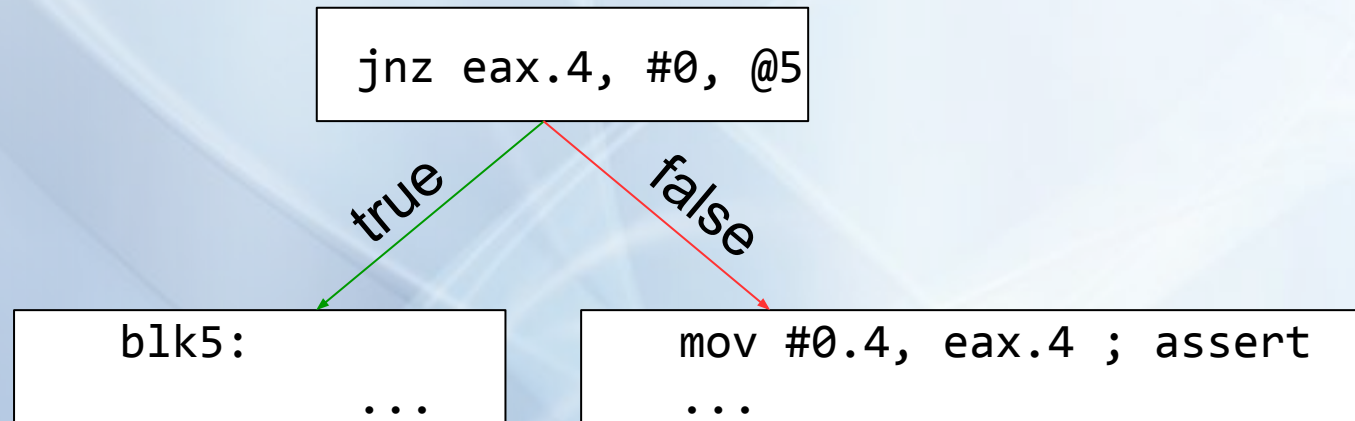
- The initial “preoptimization” step uses very simple constant and register propagation algorithm
- It is very fast
- It gets rid of most temporary registers and reduces the microcode size by two
- Normally we use a more sophisticated propagation algorithm
- It also works on the basic block level
- It is much slower but can:
  - handle partial registers (propagate eax into an expression that uses ah)
  - move entire instruction inside another
  - work with operands other than registers (stack and global memory, pair and scattered operands)

# Global optimization

- We build the control flow graph
- Perform data flow analysis to find where each operand is used or defined
- The use/def information is used to:
  - delete dead code (if the instruction result is not used, then we delete the instruction)
  - propagate operands and instructions across block boundaries
  - generate assertions for future optimizations (we know that `eax` is zero at the target of `“jz eax”` if there are no other predecessors; so we generate `“mov 0, eax”`)

# Synthetic assertion instructions

- If jump is not taken, then we know that `eax` is zero



- Assertions can be propagated and lead to more simplifications

## Simple algebraic transformations

- We have implemented (in plain C++) hundreds of very small optimization rules. For example:

$$\begin{array}{ll} (x-y)+y & \Rightarrow x \\ x- \sim y & \Rightarrow x+y+1 \\ x^*m-x^*n & \Rightarrow x^*(m-n) \\ (x<<n)-x & \Rightarrow (2^{**n}-1)*x \\ -(x-y) & \Rightarrow y-x \\ (\sim x) < 0 & \Rightarrow x >= 0 \\ (-x)^*n & \Rightarrow x^*-n \end{array}$$

- They are simple and sound
- They apply to all cases without exceptions
- Overall the decompiler uses sound rules
- They do not depend on the compiler

## More complex rules

- For example, this rule recognizes 64-bit subtractions:

CMB18 (combination rule #18):

```
sub xlow.4, ylow.4, rlow.4
```

```
sub xhigh.4, (xdu.4((xlow.4 <u ylow.4))+yhigh.4), rhigh.4
```

=>

```
sub x.8, y.8, r.8
```

if yhigh is zero, then it can be optimized away

a special case when xh is zero:

```
sub    x1, y1, r1
```

```
neg    (xdu(1not(x1 >=u y1))+yh), rh
```

- We have a swarm of rules like this. They work like little ants :)



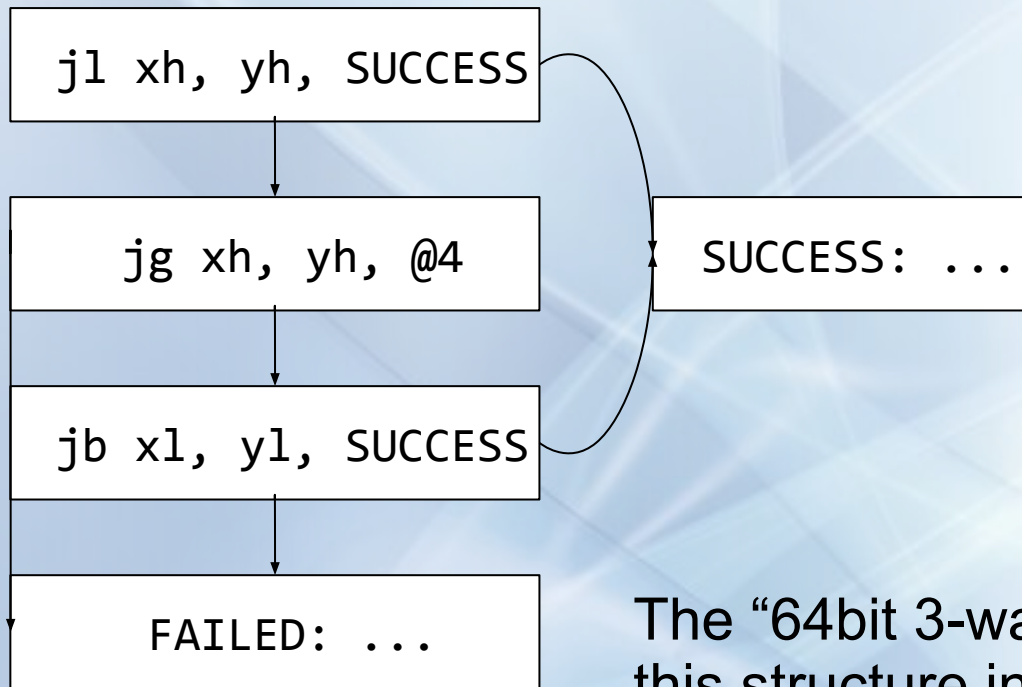
## Data dependency dependent rules

- Naturally, all these rules are compiler-independent, they use common algebraic number properties
- Unfortunately we do not have a language to describe these rules, so we manually added these rules in C++
- However, the pattern recognition does not naively check if the previous or next instruction is the expected one. We use data dependencies to find the instructions that form the pattern
- For example, the rule CMB43 looks for the 'low' instruction by searching forward for an instruction that accesses the 'x' operand:

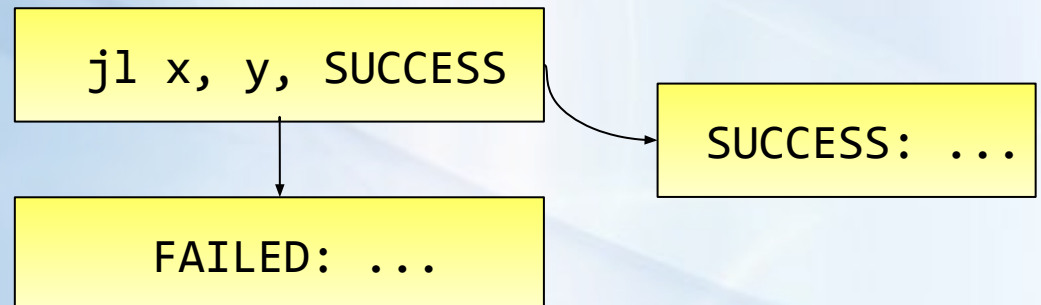
```
CMB43:  
mul    #(1<<N).4, x1.4, y1.4  
low    (x.8 >>a #M.1), yh.4, M == 32-N  
  
=>  
  
mul x.8, #(1<<N).8, y.8
```

# Interblock rules

- Some rules work across multiple blocks:



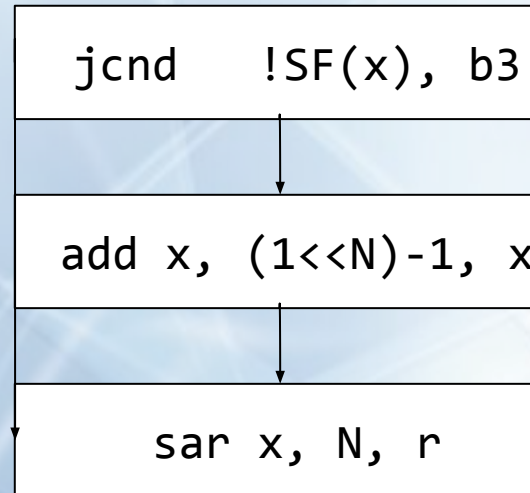
The “64bit 3-way check” rule transforms this structure into simple:



*(xh means high half of x  
xl means low half of x  
yh means high half of y  
yl means low half of y)*

## Interblock rules: signed division by power2

- Signed division is sometimes replaced by a shift:



A simple rule transforms it back:

```
sdiv x, (1<<N), r
```

# Hooks

- It is possible to hook to the optimization engine and add your own transformation rules
- The Decompiler SDK has some examples how to do it
- Currently it is not possible to disable an existing rule
- However, since (almost?) all of them are sound and do not use heuristics, it is not a problem
- In fact the processor specific parts of the decompiler internally use these hooks as well

## ARM hooks

- For example, the ARM decompiler has the following rule:

```
ijmp cs, initial_lr => ret
```

so that a construct like this: BX LR

will be converted into: RET

only if we can prove that the value of LR at the "BX LR" instruction is equal to the initial value of LR at the entry point.

- However, how do we find if we jump to the initial\_lr? Data analysis is to help us

# Data flow analysis

- In fact virtually all transformation rules are based on data flow analysis. Very rarely we check the previous or the next instruction for pattern matching
- Instead, we calculate the use/def lists for the instruction and search for the instructions that access them
- We keep track of what is used and what is defined by every microinstruction (in red). These lists are calculated when necessary:

```
mov    %argv.4, ebx.4      ; 4014E9 u=arg+4.4    d=ebx.4
mov    %argc.4, edi.4     ; 4014EC u=arg+0.4    d=edi.4
mov    &($dword_41D128).4, ST18_4.4 ; 4014EF u=          d=ST18_4.4
goto   @12                ; 4014F6 u= d=
```

## Use-def lists

- Similar lists are maintained for each block. Instead of calculating them on request we keep them precalculated:

```
; 1WAY-BLOCK 6 INBOUNDS: 5 OUTBOUNDS: 58 [START=401515 END=401517]
; USE: ebx.4, ds.2, (GLBLOW, GLBHIGH)
; DEF: eax.4, (cf.1, zf.1, sf.1, of.1, pf.1, edx.4, ecx.4, fps.2, fl.1,
;         c0.1, c2.1, c3.1, df.1, if.1, ST00_12.12, GLBLOW, GLBHIGH)
; DNU: eax.4
```

- We keep both “must” and “may” access lists
- The values in parenthesis are part of the “may” list
- For example, an indirect memory access may read any memory:

```
add    [ds.2:(ebx.4+#4.4)].4, #2.4, ST18_4.4
; u=ebx.4, ds.2, (GLBLOW, GLBHIGH)
; d=ST18_4.4
```

## Usefulness of use-def lists

- Based on use-def lists of each block the decompiler can build global use-def chains and answer questions like:
  - Is a defined value used anywhere? If yes, where exactly? Just one location? If yes, what about moving the definition there? If the value is used nowhere, what about deleting it?
  - Where does a value come from? If only from one location, can we propagate (or even move) it?
  - What are the values that are used but never defined? These are the candidates for input arguments
  - What are the values that are defined but never used but reach the last block? These are the candidates for the return values



# Global propagation in action

- Image we have code like this:

blk1

```
mov #5.4, esi.4
```

blk2

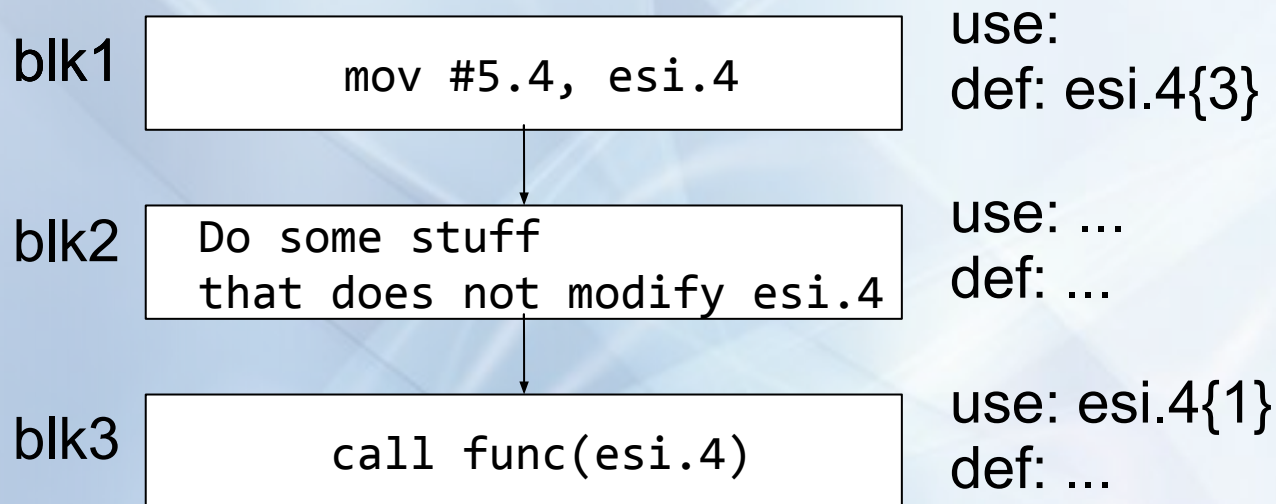
```
Do some stuff  
that does not modify esi.4
```

blk3

```
call func(esi.4)
```

# Global propagation in action

- The use-def chains clearly show that esi is defined only in block #1:



Therefore it can be propagated:

```
call func(#5.4)
```

# Data flow analysis

- The devil is in details
- Our analysis engine can handle partial registers (they are a pain)
- Big endian and little endian can be handled as well (however, we sometimes end up with the situations when a part of the operand is little endian and another part – big endian)
- The stack frame and registers are handled
- Registers can be addressed only directly
- Stack location can be addressed indirectly and our analysis takes this into account
- Well, we have to make some assumptions...

# Aliasability

- Take this example:

```
mov #1.4, %stkvar      ; store 1 into stkvar
stx #0.4, ds.2, eax.4  ; store 0 into [eax]
call func(%stkvar)
```

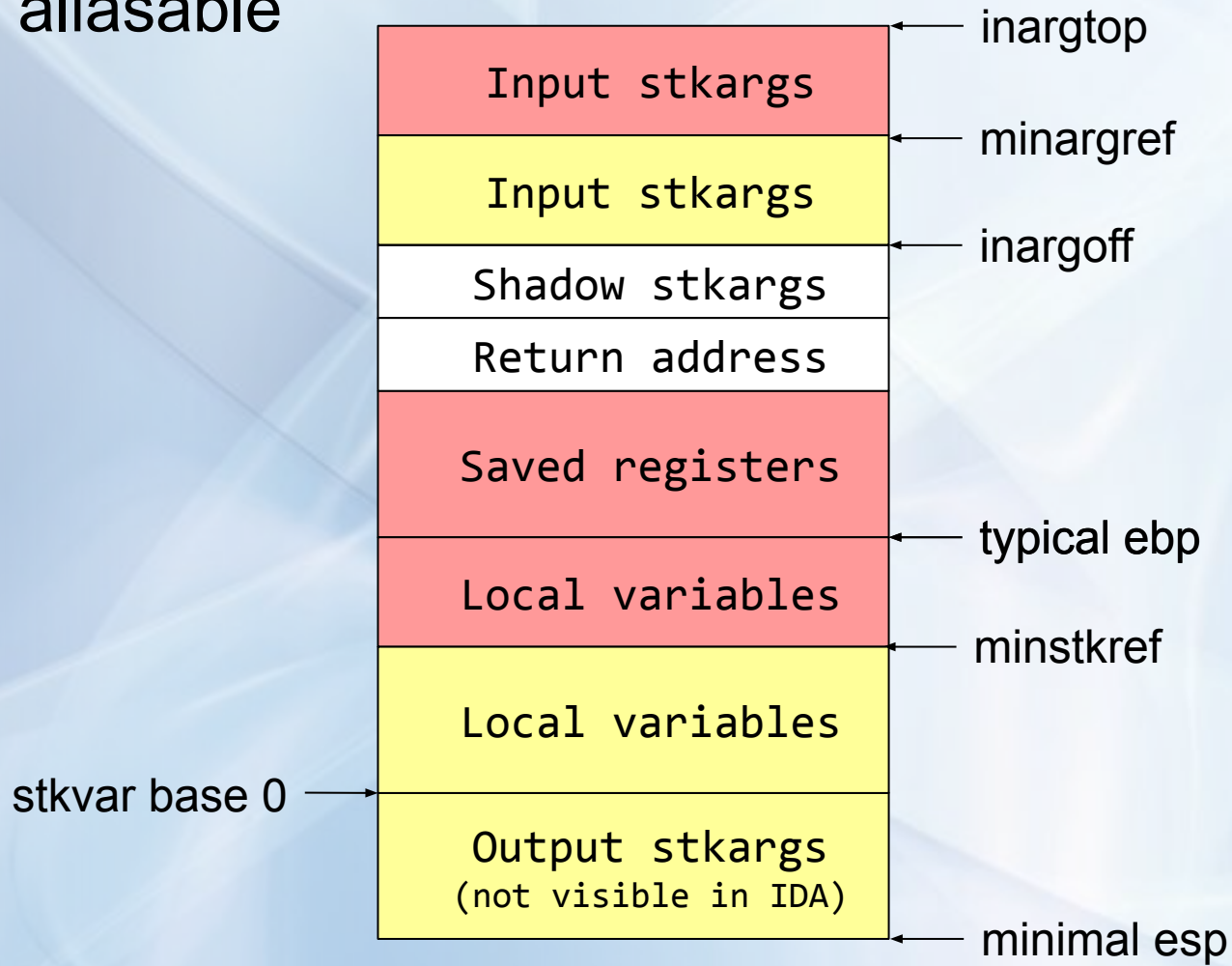
can we claim that `%stkvar == 1` after `stx`?

- Naturally, in general case we can not
- But it turns out that in some case we can claim it
- Namely:
  - If we haven't taken the address of any stack variable
  - Or, if we did, the address we took is higher (\*)
  - Or, if the address is lower, it was not moved into `eax`
- Overall it is a tough question

(\*)note: yes, this is one of the assumptions our decompiler makes

# Stack as viewed by the decompiler

- Yellow part is mapped to microregisters
- Red is aliasable



## Minimal stack reference

- Aliasability is unsolvable problem in general
- We should optimize things only if we can prove the correctness of the transformation
- We keep track of expressions like `&stkvar` and calculate the minimal reference (`minstkref`)
- We assume that everything below `minstkref` can be accessed only directly, i.e. is not aliasable
- We propagate this information over the control graph
- One value is maintained per block (we could probably improve things by calculating `minstkref` for each instruction)
- A similar value is maintained for the incoming stack arguments (`minargref`)

# Minstkref propagation

- We use the control flow graph:

```
lea ecx, [esp+10] ; take offset 10
call func          ; probably uses ecx
mov rax, [esp+14] ; stkvar sp+14
...
```

minstkref=10

```
lea ecx, [esp+20] ; take offset 20
call func          ; probably uses ecx
mov rax, [esp+14] ; microregister ST14
...
```

minstkref=20

```
mov rax, [esp+14] ; stkvar sp+14
...
```

minstkref=10

## Testing the microcode

- Microcode is verified for consistency after every transformation
- BTW, third party plugins should do the same
- Very few microcode related bug reports
- We have quite extensive test suites that constantly grow
- A hundred or so of processors cores running tests
- However, after publishing microcode there will be a new wave of bug reports
- Found a bug? Send us the database with the description how to reproduce it
- Most problems are solved within one day or faster



## Publishing microcode

- The microcode API for C++ will be available in the next version of IDA
- Python API won't be available yet
- We will start beta testing the next week
- Decompiler users with active support: feel free to send an email to [support@hex-rays.com](mailto:support@hex-rays.com) if you want to participate
- Check out the sample plugins that show how to use the new API

**Was it interesting?**

**Thank you for your attention!  
Questions?**