

ALGORITHMS AND DATA STRUCTURES

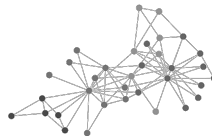
LECTURE 6 - SORTING

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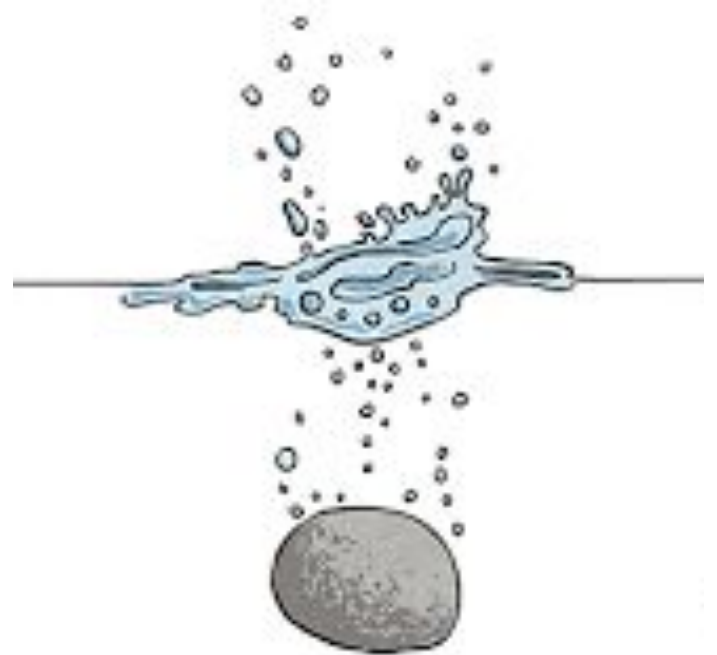
CONTENT

1. Bubble sort
2. Heap Sort
3. Divide and Conquer
4. Merge Sort
5. Quick Sort



Bubble Sort – $O(n^2)$

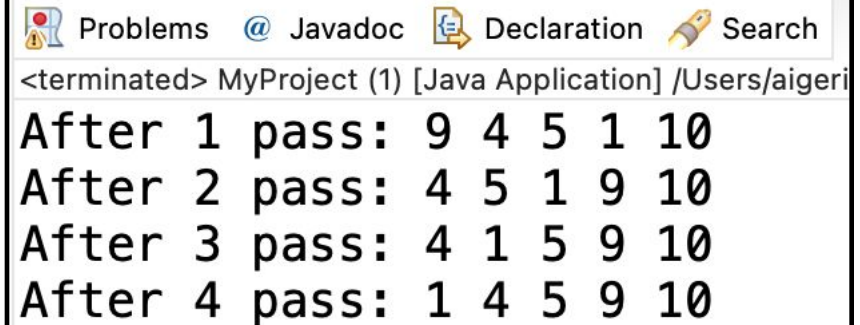
- Bubble Sort is the simplest sorting algorithm
- Several passes through the array
- Successive pairs of elements are compared
- Repeatedly swaps the adjacent elements if they are in wrong order
- At each i 'th iteration of the outer loop the maximum (can be minimum) element is moved to the position of $n-i-1$



Bubble Sort – $O(n^2)$

```
public static int[] bubbleSort(int[] arr) {  
    for(int i = 0; i < arr.length - 1; i++) {  
        for(int j = 0; j < arr.length - 1 - i; j++) {  
            if(arr[j] > arr[j+1]) {  
                int temp = arr[j];  
                arr[j] = arr[j+1];  
                arr[j+1] = temp;  
            }  
        }  
    }  
    return arr;  
}  
  
public static void main(String[] args) {  
    int[] arr = {10, 9, 4, 5, 1};  
    arr = bubbleSort(arr);  
    for(int i = 0; i < arr.length; i++)  
        System.out.println(arr[i]);  
}
```

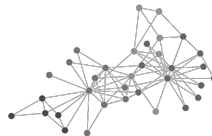
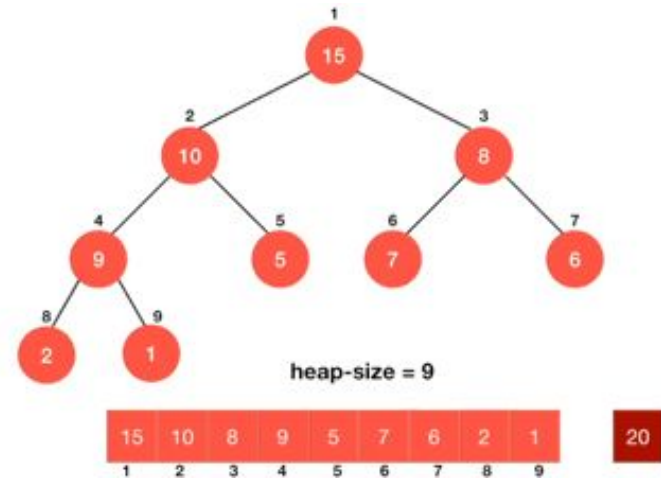
Output:



Problems @ Javadoc Declaration Search
<terminated> MyProject (1) [Java Application] /Users/aigeri
After 1 pass: 9 4 5 1 10
After 2 pass: 4 5 1 9 10
After 3 pass: 4 1 5 9 10
After 4 pass: 1 4 5 9 10

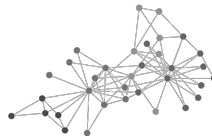
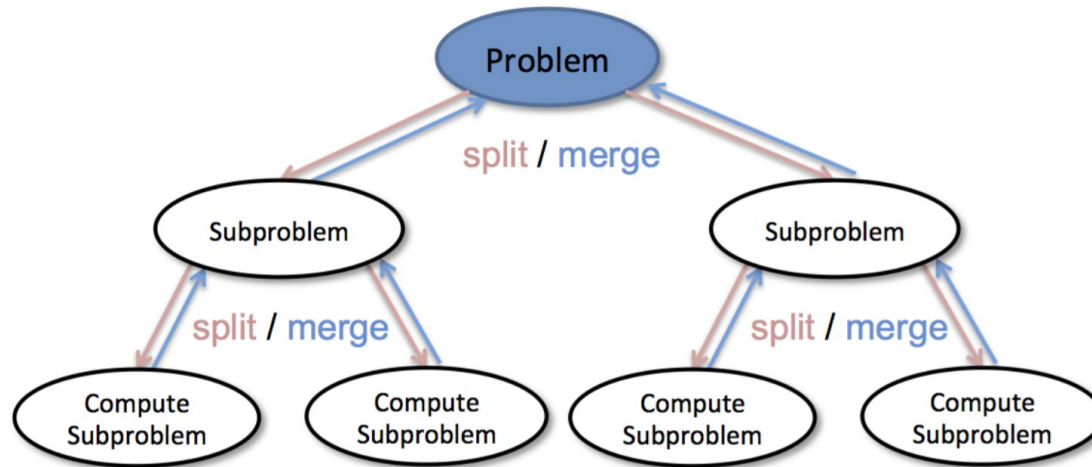
HEAP SORT – $O(N \log N)$

1. Heap tree must be built from the input data to implement heap sort
2. Swap the root with the last item of the heap followed by reducing the size of heap by 1 (**removing** the root element)
3. **Heapify** the root of the tree
4. Repeat step 2 and 3 while size of heap is greater than 1



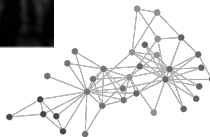
DIVIDE-AND-CONQUER

Divide and conquer algorithm works by recursively breaking down a problem into two or more sub-problems of the same type, until these become simple enough to be solved directly



JOHN VON NEUMANN

- Invented a Merge Sort algorithm in 1945, in which the first and second halves of an array are each sorted recursively and then merged



MERGE SORT

Divide data sequence recursively till each data sub-sequence obtains its "**atomic**" representation

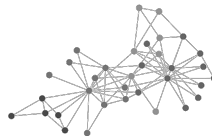
The sequence is said to be sorted if data is arranged in an ordered way OR **the length of sequence is exactly 1**

Merge Sort uses $\leq N \lg N$ compares to sort an array of length N



```
5 3 7 1 0 8 5
```

```
mergesort([5, 3, 7, 1, 0, 8, 5])
```



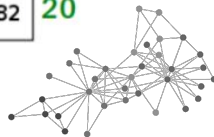
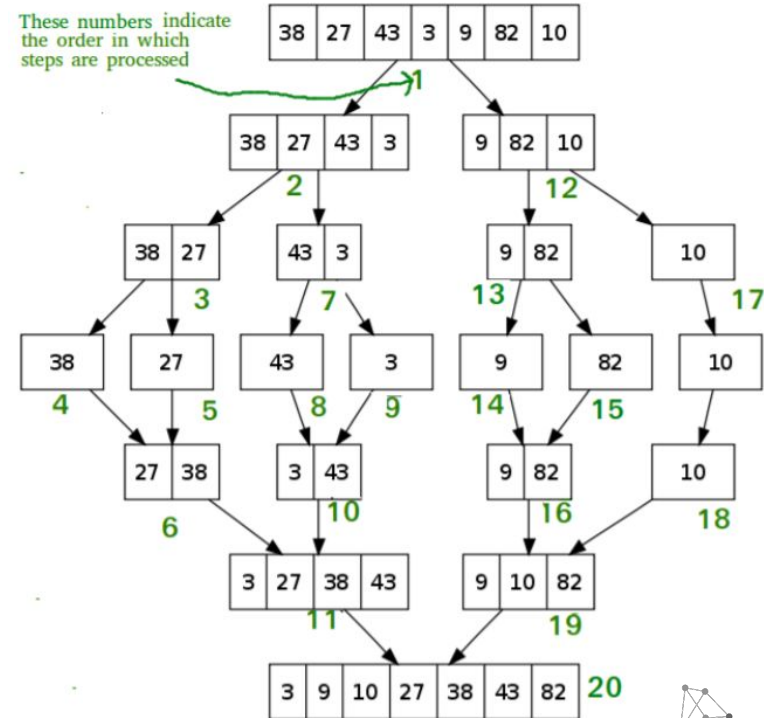
MERGE SORT (CONTINUED)

The initial array length is 7

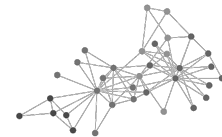
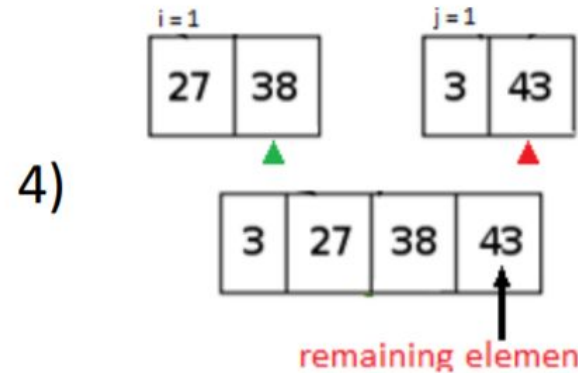
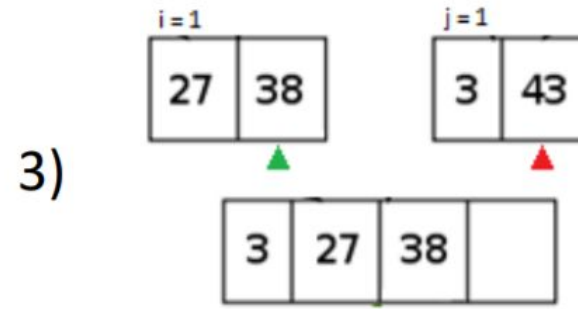
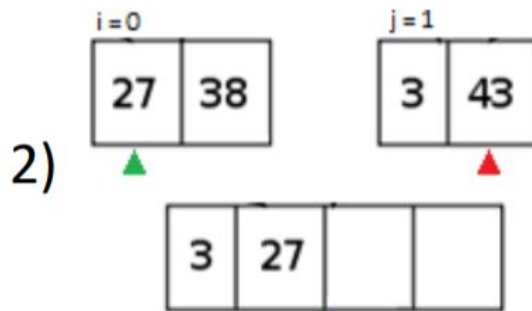
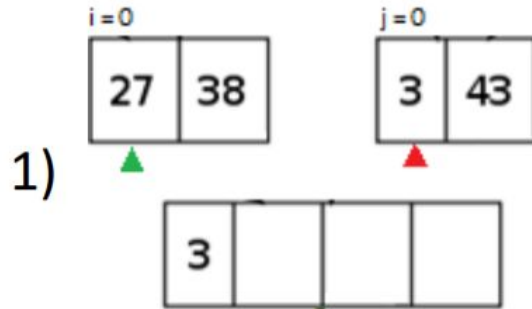
It is divided to two subarrays with lengths 4 and 3

Division is repeated recursively till all subarrays reach an atomic view (one element in each array)

Since array that contains only one element is said to be sorted, merging process is started



MERGING TWO SORTED SUBARRAYS INTO ONE



SAMPLE CODE TO START

```
public static void sort(int[] x) {
    sort(x, 0, x.length-1);
}

private static void sort(int[] x, int start, int end) {
    if (start < end) {
        int middle = ???;
        sort(x, ???, ???);
        sort(x, ???, ???);

        merge(x, start, end, middle);
    }
}

public static void merge(int[] x, int start, int end, int middle) {

    int[] a = new int[???];
    int[] b = new int[???];

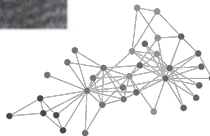
    // take a copy of a left half
    for (int j = 0; j <= ???; j++) {
        a[j] = x[???];
    }

    // take a copy of a right half
    for (int j = 0; j <= ???; j++) {
        b[j] = x[???];
    }
}
```



SIR CHARLES ANTONY RICHARD HOARE

- Invented a Quick Sort algorithm in 1959/60. His sorting method is based on divide-and conquer algorithm



QUICK SORT

Select a pivot point which could be:

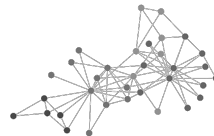
- First element
- Last element
- Middle element in the sequence

5	4	2	1	3
---	---	---	---	---

Partition the other elements into two sub-arrays, according to whether they are less or greater than the pivot point

Generally Quick Sort uses $\leq O(N \log N)$ compares to sort an array of length N

Worst case: $O(N^2)$ – when the selected pivot point is always the smallest or largest



QUICK SORT (CONTINUED)

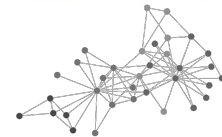
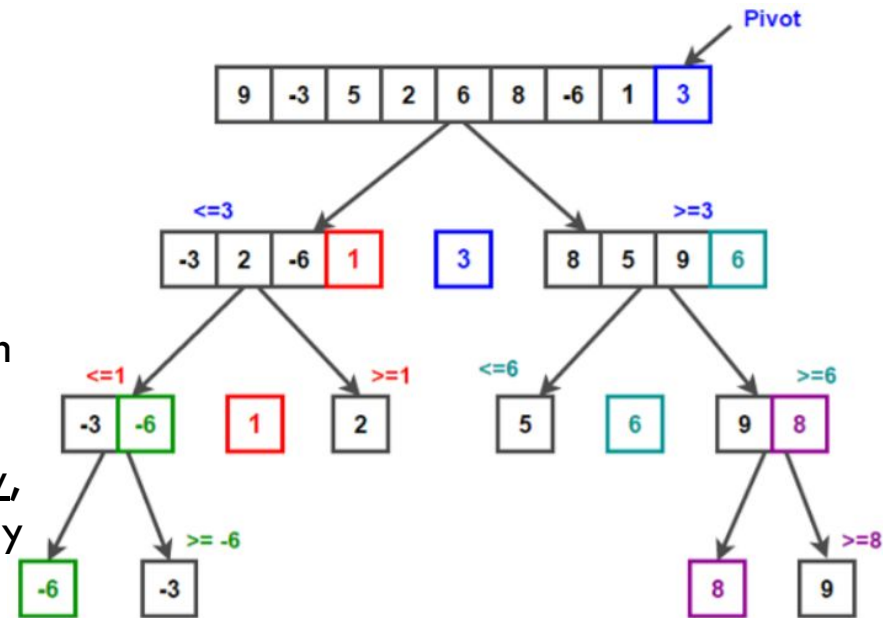
Select a pivot point (e.i. last element)

Partition to two sub-arrays (using indices)

- Less elements to the left
- Greater elements to the right

Repeat recursively till all subarrays reach an atomic view (one element in each array)

No need to create two sub-arrays physically,
since swaps are applied to the original array
by **manipulating indices**



SAMPLE CODE TO START

```

public void quickSort(int[] x) {
    quickSort(x, 0, x.length - 1);
}

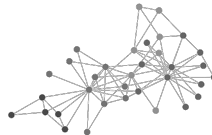
private void quickSort(int[] x, int start, int end)
{
    if (start < end)
    {
        int pi = partition(x, start, end);
        quickSort(x, ???, ???); // sort left half
        quickSort(x, ???, ???); // sort right half
    }
}

private int partition(int[] x, int start, int end)
{
    int pivot = x[???]; // select pivot point

    // Your code goes here

    return ???;
}

```



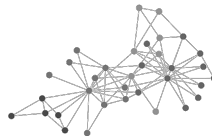
LITERATURE

Algorithms, 4th Edition, by Robert Sedgwick and Kevin Wayne, Addison-Wesley

□ Chapter 2.1-2.3

Grokking Algorithms, by Aditya Y. Bhargava, Manning

□ Chapter 4



**GOOD
LUCK!**

