



ALGORITHMS AND DATA STRUCTURES

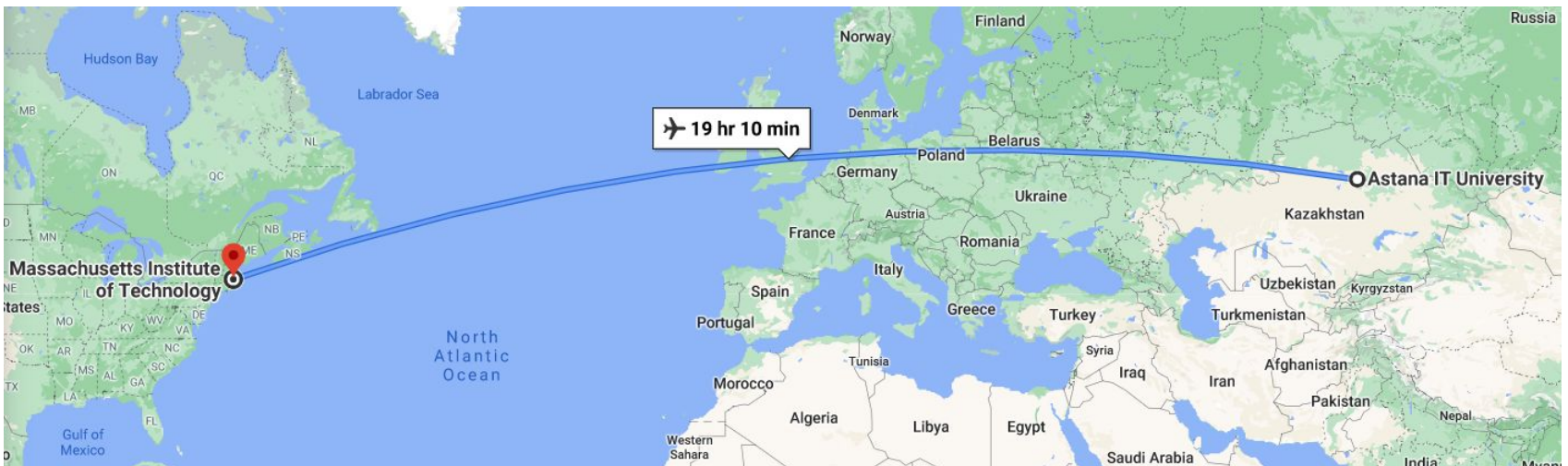
Introduction and Overview

LECTURE #1

◦ Yerasyl Amanbek

Define Algorithm (1)

- **Origin:**
 - Al-Khwārizmī, rendered as (Latin) *Algoritmi*
- **Definition:**
 - Procedure approach to solve “computational problems”
- **Example:**
 - Find the shortest path from AITU to MIT

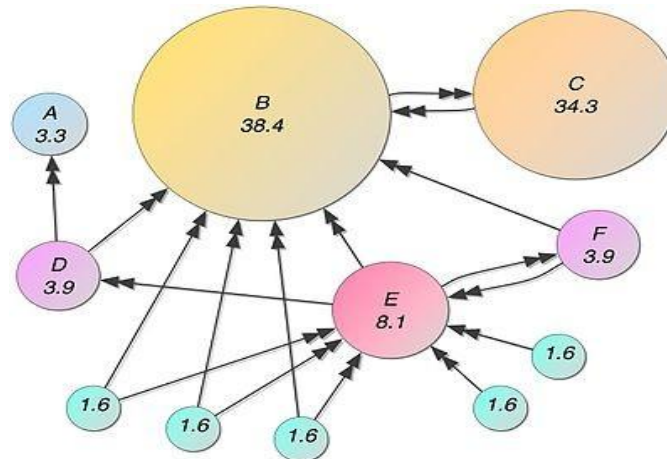


Define Algorithm (2)

A *procedure* for solving a problem in terms of

1. the **actions** to execute and
2. the **order** in which these actions execute

is called an **algorithm**. The following example demonstrates that correctly specifying the order in which the actions execute is important.



Find the most relevant web page for a query (Google) : PageRank

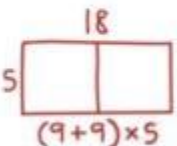
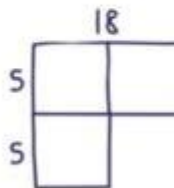
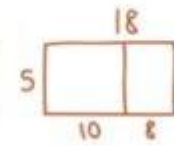
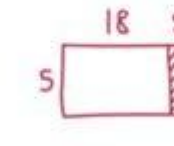
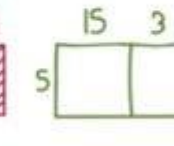
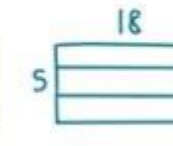
Main Features of an Algorithm

- Various features
 - Reusability/modularity
 - Simplicity
 - Memory footprint
 - Speed
- Algorithm is all about efficiency: **Time vs. Space**
- **Time complexity:** Developing a formula for predicting *how fast* and algorithm is, based on input size.
- **Space complexity:** Developing a formula for predicting *how much memory* an algorithm requires, based on input size.
- Memory is extensible, time is not!

Data Structures and Algorithms (1)

- To solve a given problem by using computers, you need to design an algorithm for it.
- Multiple algorithms can be designed to solve a problem.

Work out 18×5 and show a visual solution.

Neil	Ricardo	Sammi	Jaime	Ariane	Bryan
					
$(9+9) \times 5$	$18 \times 5 = 9 \times 10$	$(10 \times 5) + (8 \times 5)$	$20 \times 5 = 100$ $2 \times 5 = 10$ $100 - 10 = 90$	$15 \times 5 = 75$ $3 \times 5 = 15$ $75 + 15 = 90$	$(18 \times 2) + (18 \times 3) + 18$ $36 + 36 + 18 = 90$

- An algorithm that provides the maximum efficiency should be used for solving the problem.

Data Structures and Algorithms (2)

- The efficiency of an algorithm can be improved by using an appropriate data structure.
- Data structures help in creating programs that are simple, reusable, and easy to maintain.
- To solve a problem, you need a computer to write a program.
- A program is made up of two parts:
 - Algorithm
 - Data structures
 - Arrays
 - Queues
 - Lists
 - Linked Lists
 - Trees
 - Graphs

Example of an Algorithm

Consider the “rise-and-shine algorithm” followed by one executive for getting out of bed and going to work.

- (1) Get out of bed;
- (2) take off pajamas;
- (3) take a shower;
- (4) get dressed;
- (5) eat breakfast;
- (6) carpool to work.

This routine gets the executive to work well prepared to make critical decisions. Suppose that the same steps are performed in a slightly different order:

- (1) Get out of bed;
- (2) eat breakfast;
- (3) take off pajamas;
- (4) take a shower;
- (5) Get dressed;
- (6) carpool to work.

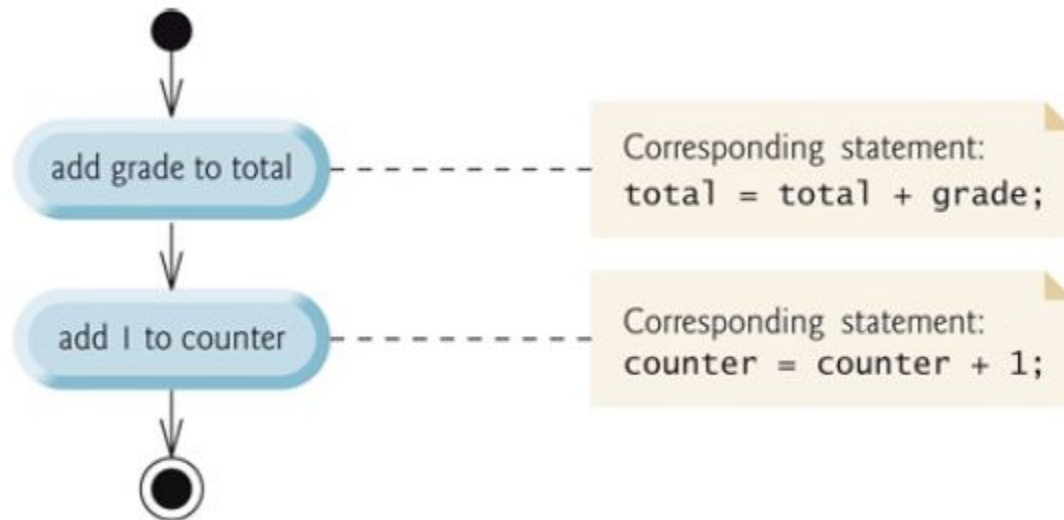
Pseudocode

- Informal language that helps to understand and develop algorithms
- Pseudocode is similar to everyday language.
- Pseudocode:
 - Does not execute on computers
 - Help to “think out”
 - Can be easily converted to program

Control

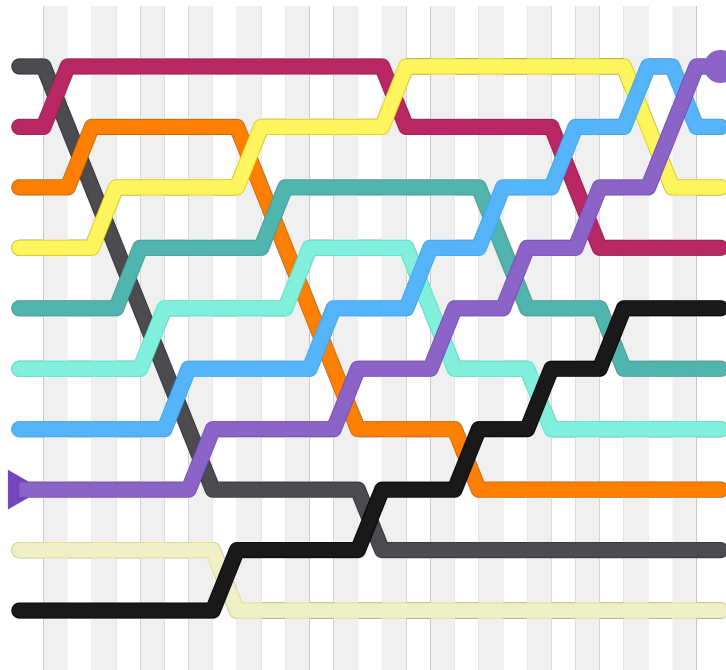
Structures

- Sequential execution – execution of statements in the order in which they are written
- Activity diagram:



Bubble sort algorithm

- Bubble Sort is the simplest sorting algorithm
- Several passes through the array
- Successive pairs of elements are compared
- Repeatedly swaps the adjacent elements if they are in wrong order
- At each i 'th iteration of the outer loop the maximum (can be minimum) element is moved to the position of $n-i-1$



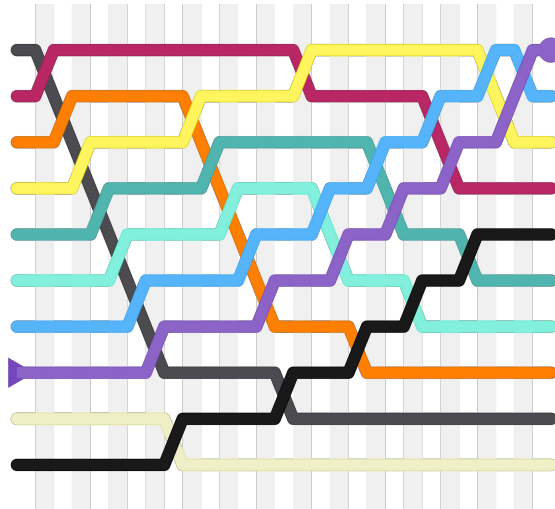
Bubble sort algorithm

Pseudocode:

Input: An array of n numbers, $A[1\dots n]$

Bubble-Sort(A)

1. for $i = (A.length-1)$ to 0
2. for $j = 0$ to $(i - 1)$
3. if $(A[j] > A[j+1])$
4. swap $A[j]$ and $A[j+1]$



Reverse Array

```
n = 5;
```

```
A[n] = {1, 2, 3, 4, 5};
```

```
for (i = 0; i < n/2; i++){
```

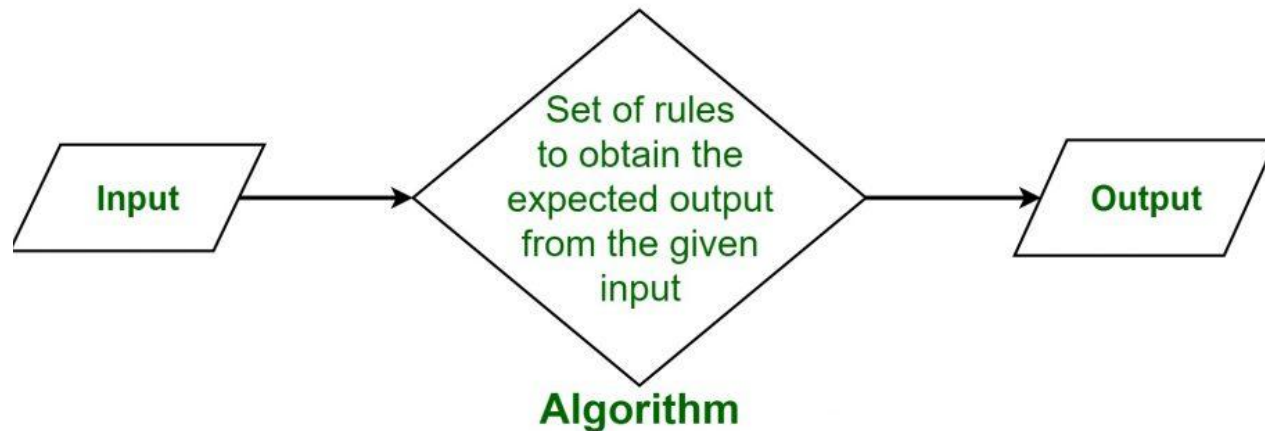
```
    temp = A[i];
```

```
    A[i] = A[n - i - 1];
```

```
    A[n - i - 1] = temp;
```

```
}
```

- Here we do not need to initialize second array in order to reverse on array.
- Reversing of the array also can be considered as an algorithm, as an algorithm can be defined as set of rules to obtain expected output



Checking array

- In order to check array: is it sorted or not?

```
n = 5;
```

```
A[n] = {1, 2, 3, 5, 4};
```

```
bool is_sorted = true;
```

```
for (i = 0; i < n-1; i++){
```

```
//here we check by pair so, we will not check n and n+1 element
```

```
    if (A[i] > A[i+1]){
```

```
        is_sorted = false;
```

```
        break;
```

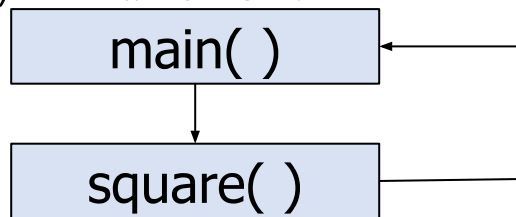
```
    }
```

```
}
```

Recursion

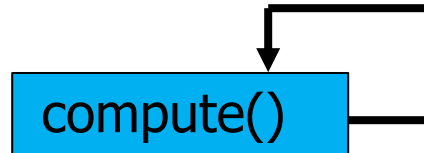
- So far, we have seen methods that call other functions.

– For example, the `main()` method calls the `square()` function.



- Recursive Method:

– A recursive method is a method that calls itself.



Why we need Recursion?

- Some problems are more easily solved by using recursive functions.
- If you go on to take a computer science algorithms course, you will see lots of examples of this.
- For example:
 - Traversing through a directory or file system.
 - Traversing through a tree of search results.
- For today, we will focus on the basic structure of using recursive methods.

Simplest recursion code:

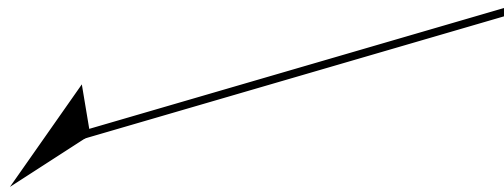
```
#include <iostream>
using namespace std;
void count(int);
int main()
{
    count(0);
    cout<<endl;
}
```

```
void count (int index)
{
    cout<<index<<" ";
    if (index < 2) {
count(index+1);
    }
}
```

This program simply counts from 0-2:

012

This is where the recursion occurs. You can see that the count() function calls itself.



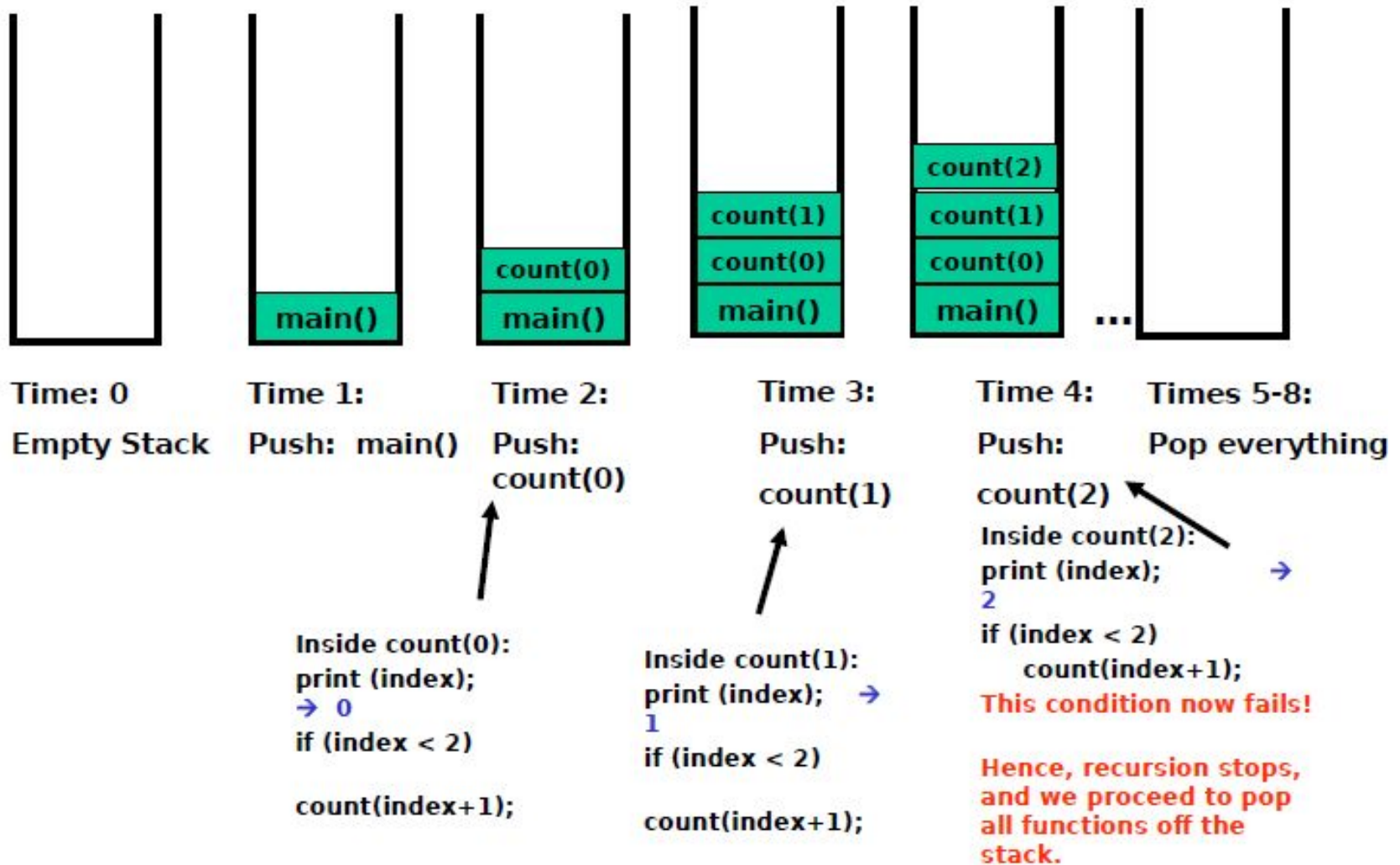
Visualizing recursion

- To understand how recursion works, it helps to visualize what's going on.
- To help visualize, we will use a common concept called the *Stack*.
- A stack basically operates like a container of trays in a cafeteria. It has only two operations:
 - Push: you can push something onto the stack.
 - Pop: you can pop something off the top of the stack.

Stacks and Methods

- When you run a program, the computer creates a stack for you.
- Each time you invoke a method, the method is placed on top of the stack.
- When the method returns or exits, the method is popped off the stack.
- The diagram on the next page shows a sample stack for a simple Java program.

Stacks and Recursion



Finding factorial:

- Computing factorials are a classic problem for examining recursion.

- A factorial is defined as follows:

- $n! = n * (n-1) * (n-2) \dots * 1;$

- For example:

$$1! = 1 \text{ (Base Case)}$$

$$2! = 2 * 1 = 2$$

$$3! = 3 * 2 * 1 = 6$$

$$4! = 4 * 3 * 2 * 1 = 24$$

$$5! = 5 * 4 * 3 * 2 * 1 = 120$$

If you study this table closely, you will start to see a pattern. The pattern is as follows:

You can compute the factorial of any number (n) by taking n and

multiplying it by the factorial of (n-1).

For example:

$$5! = 5 * 4!$$

$$\text{(which translates to } 5! = 5 * 24 = 120)$$

Seeing the Pattern

- Seeing the pattern in the factorial example is difficult at first.
- But, once you see the pattern, you can apply this pattern to create a recursive solution to the problem.
- Divide a problem up into:
 - What it can do (usually a base case)
 - What it cannot do
- What it cannot do resembles original problem
- The function launches a new copy of itself (recursion step) to solve what it cannot do.

Recursion vs. Iteration

- Iteration

- Uses repetition structures (for, while or do...while)
- Repetition through explicitly use of repetition structure
- Terminates when loop-continuation condition fails
- Controls repetition by using a counter

- Recursion

- Uses selection structures (if, if...else or switch)
- Repetition through repeated method calls
- Terminates when base case is satisfied
- Controls repetition by dividing problem into simpler one

Recursion vs. Iteration

(cont.)

- Recursion
 - More overhead than iteration
 - More memory intensive than iteration
 - Can also be solved iteratively
 - Often can be implemented with only a few lines of code